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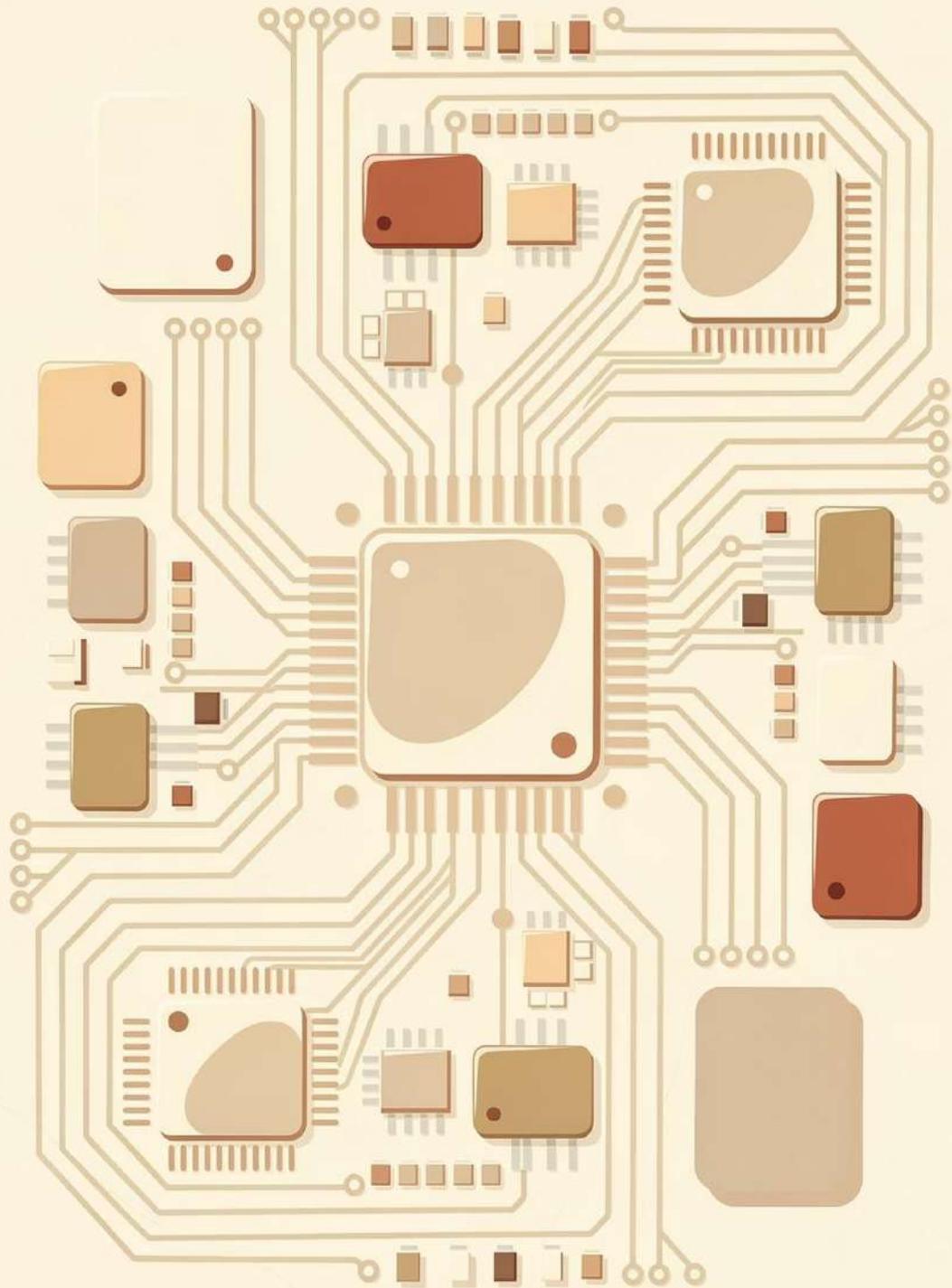
Department of Computer Science and Engineering

23CST206-OPERATING SYSTEMS AND VIRTUALIZATION

B.E- CSE /IV SEMESTER

UNIT - II PROCESS MANAGEMENT

Topic 9: Deadlock Prevention & Avoidance



Deadlock Prevention & Avoidance in Operating Systems

Understanding critical strategies for resource management in concurrent systems

What is Deadlock?

- A deadlock occurs when processes in a system are blocked, each waiting for resources held by others creating a circular wait with no resolution.
- The Classic Scenario

Process A holds Resource 1, needs Resource 2

Process B holds Resource 2, needs Resource 1

Neither can proceed → system frozen



Four Necessary Conditions for Deadlock



Mutual Exclusion

Resources cannot be shared—only one process uses a resource at a time



Hold and Wait

Processes hold resources while waiting for additional ones



No Preemption

Resources cannot be forcibly taken—only released voluntarily

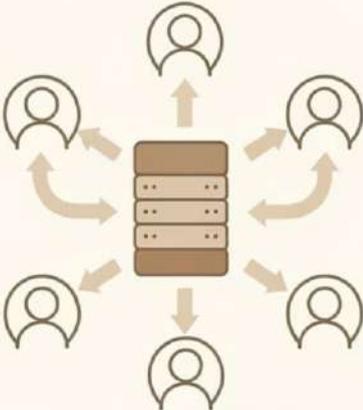


Circular Wait

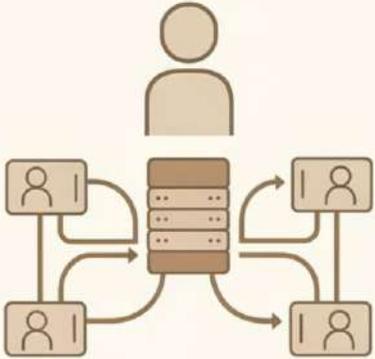
Chain of processes, each waiting for a resource held by the next

Deadlock Prevention Strategies

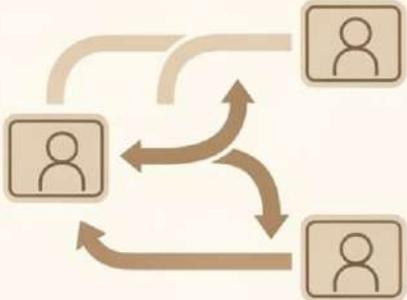
Mutual Exclusion Prevention



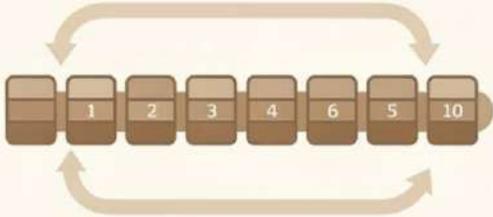
Hold and Wait Prevention



No Preemption Prevention



Circular Wait Prevention

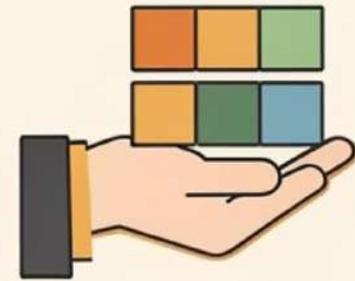


Deadlock Prevention Strategies



Break Mutual Exclusion

Make resources shareable as much as possible (e.g., read-only files)



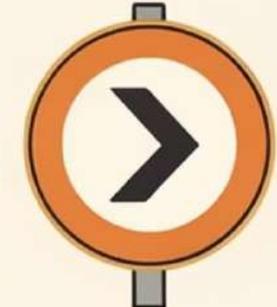
Eliminate Hold & Wait

Require processes to request all resources upfront (e.g., execution)



Eliminate Hold & Wait

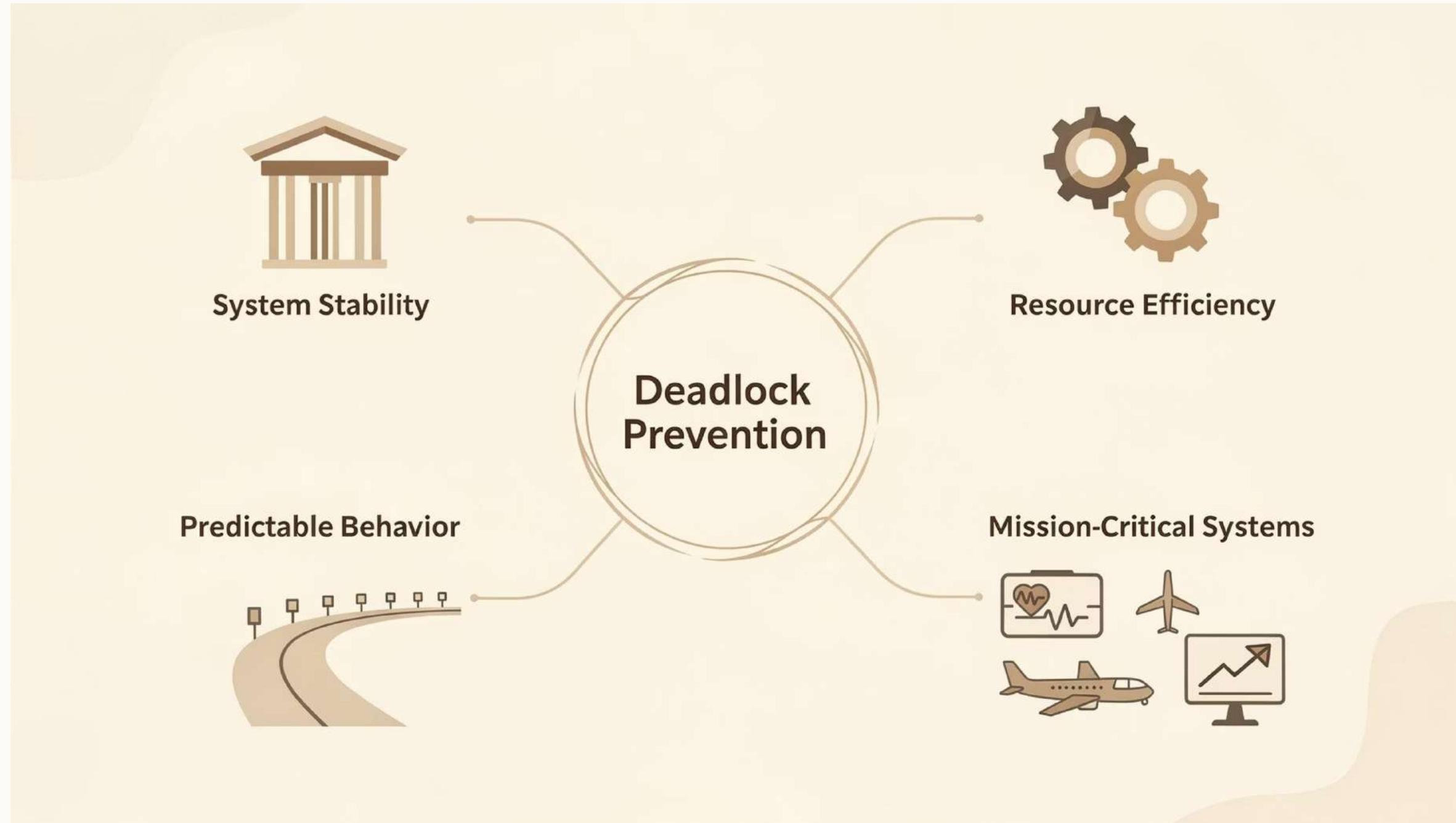
Require processes to request all resources upfront (e.g., execution)



Eliminate Hold & Wait

Require processes to request all resources upfront (e.g., execution)

Why Deadlock Prevention Matters



Deadlock Avoidance

1

Dynamic Assessment

System continuously analyzes resource allocation patterns before granting requests

2

Safe State Analysis

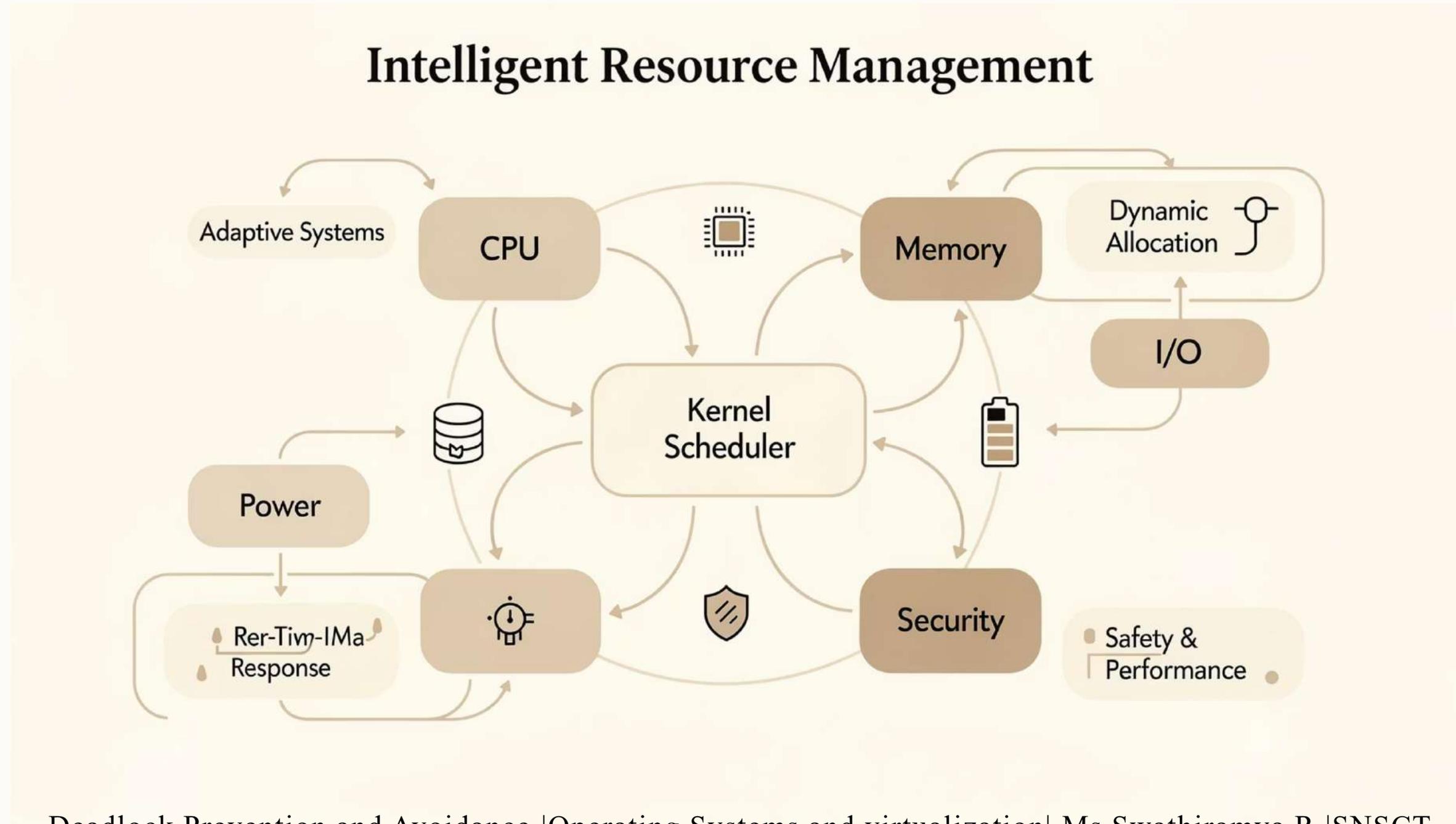
Ensures system remains in "safe state" where all processes can complete

3

Strategic Denial

May temporarily deny safe requests to prevent future unsafe states

The Role of Avoidance in Modern OS



Problem Definition

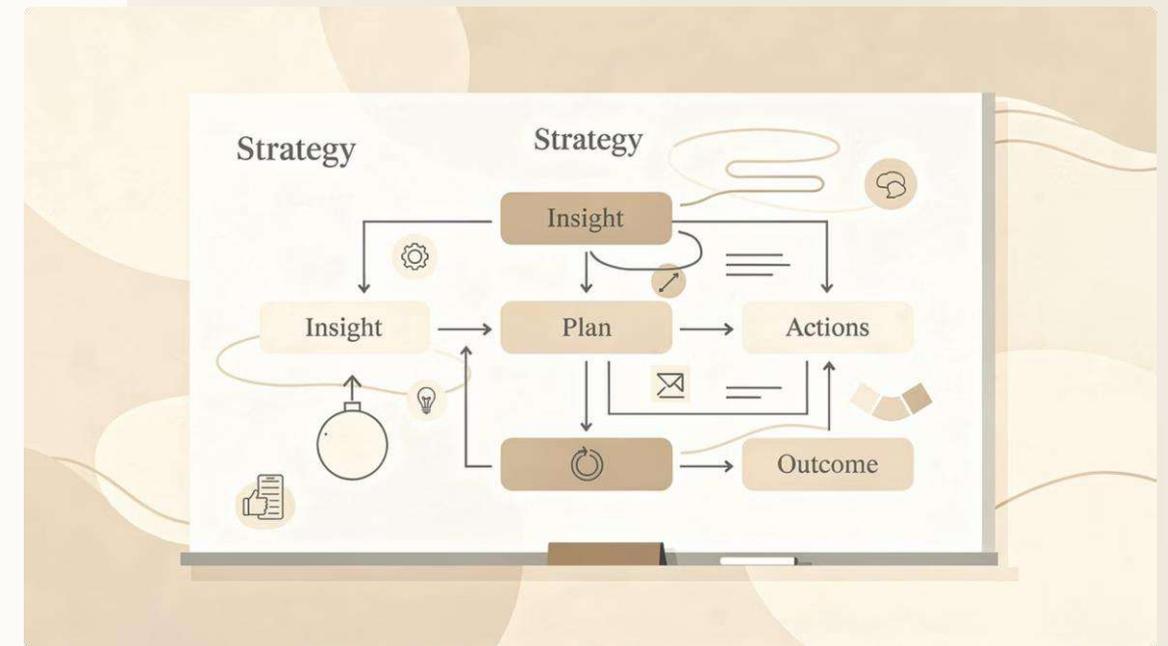
How might we manage limited system resources among competing processes to ensure continuous progress while maximizing efficiency?

Prevention Challenge

Design constraints that eliminate deadlock conditions without crippling system performance

Avoidance Challenge

Create algorithms that predict unsafe states and make real-time resource allocation decisions



Empathize: Understanding Stakeholders



System Developers

Need simple, reliable mechanisms that don't complicate code



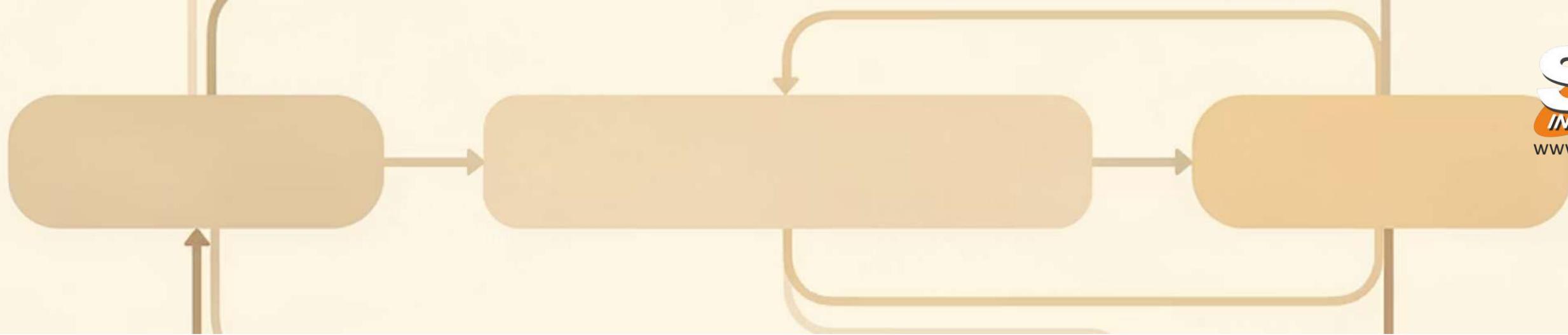
Database Admins

Require transaction integrity without sacrificing performance



End Users

Expect responsive systems that never freeze or lose data



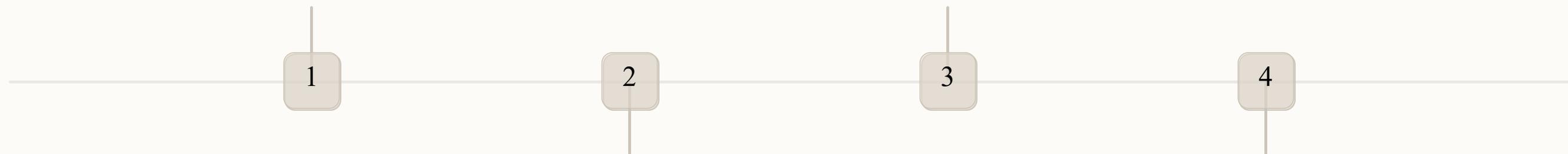
Define & Ideate: Solution Approaches

Resource Ordering

Assign global priority to resources; always request in ascending order

Wait-Die Scheme

Timestamp-based: older processes wait, younger processes die and restart



Banker's Algorithm

Track max needs; grant requests only if safe state maintained

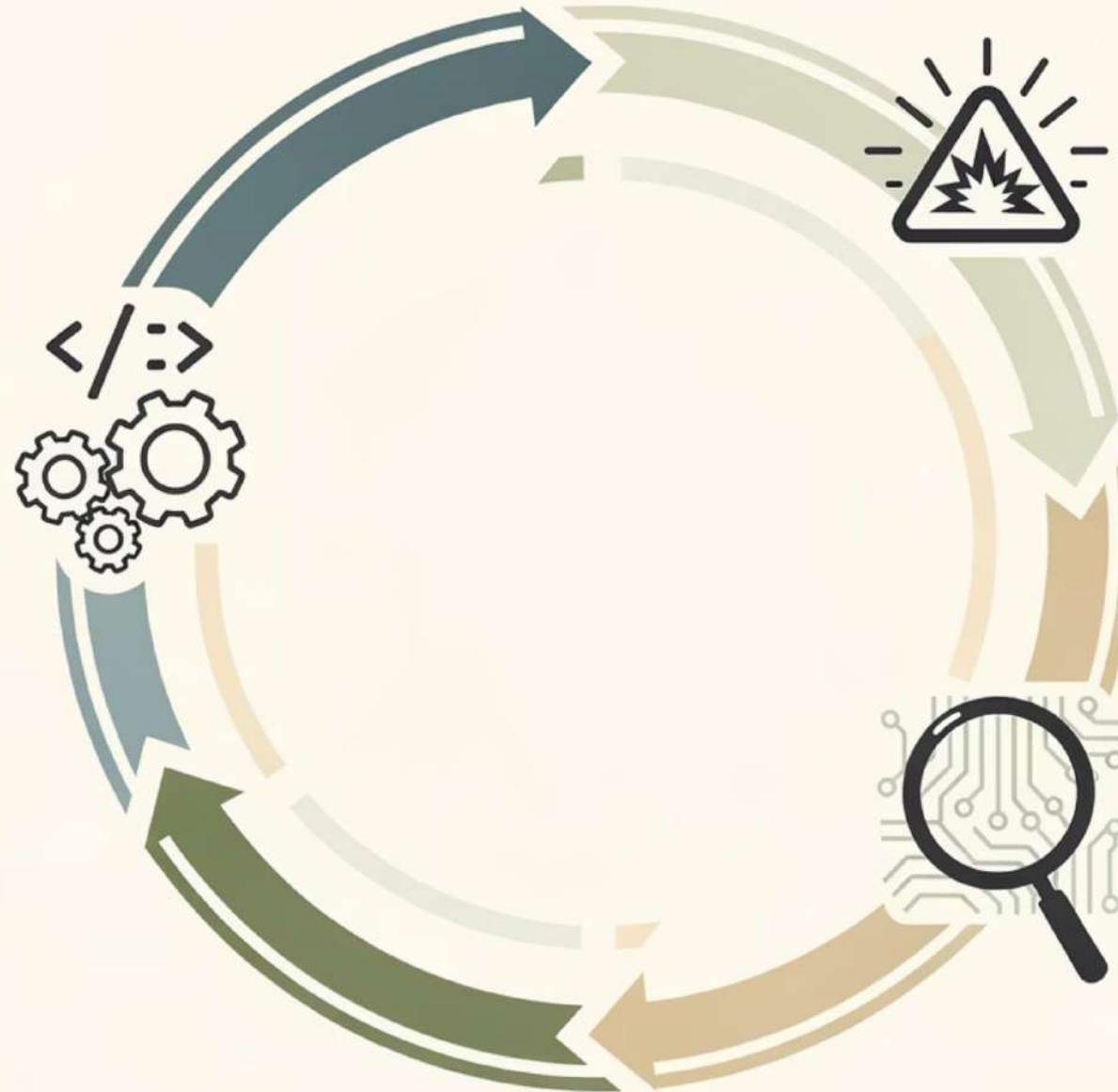
Wound-Wait Scheme

Older processes preempt younger; younger processes wait

Prototype & Test

Implementation Testing

- Simulate Workload
- Create test scenarios with varying resource demands



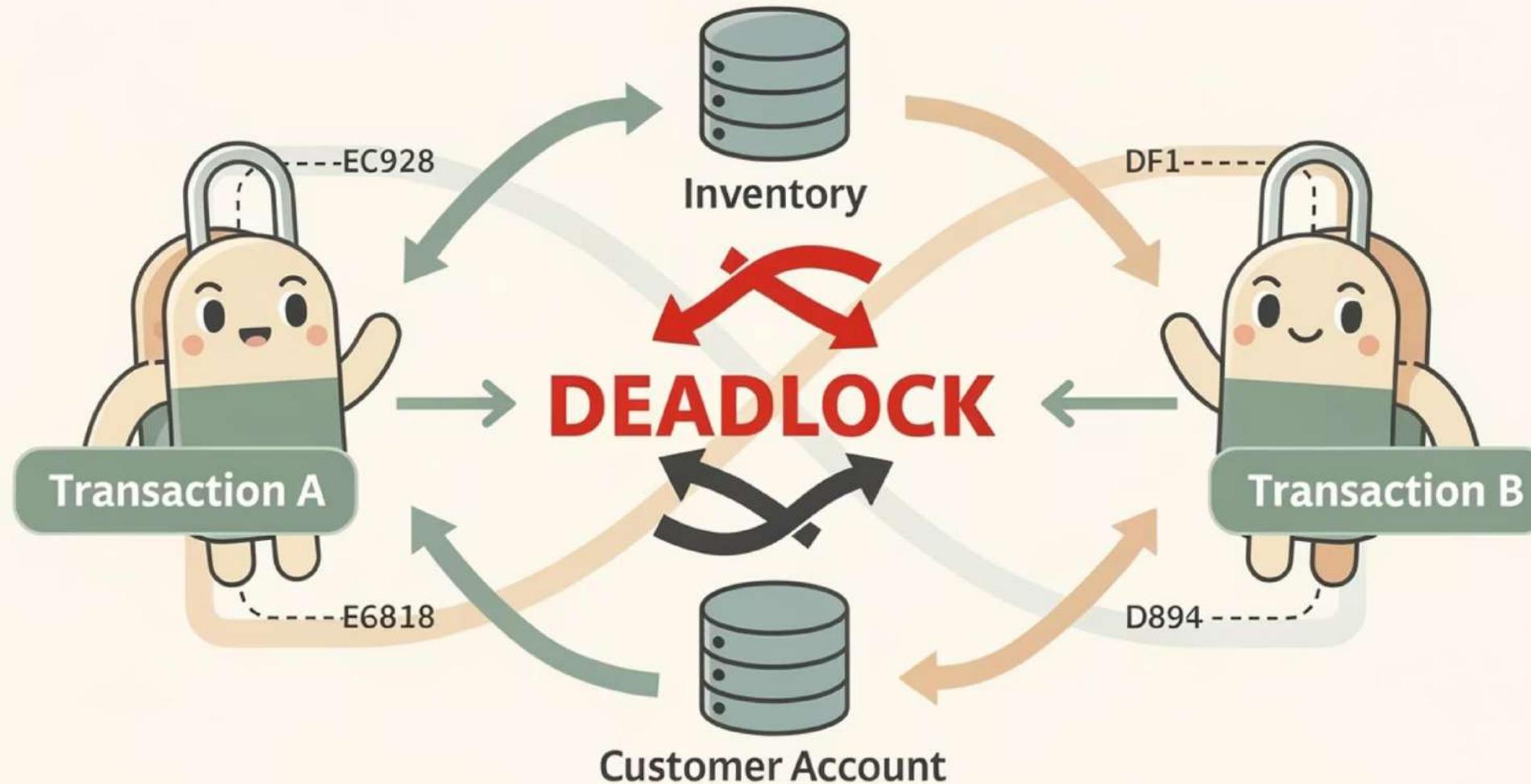
Stress Testing

- Push system to extreme conditions many concurrent processes
- Measure Performance
- Track throughput, latency, and resource utilization

Refine Algorithm

- Optimize based on real-world bottlenecks discovered

Database Transaction Management



Prevention

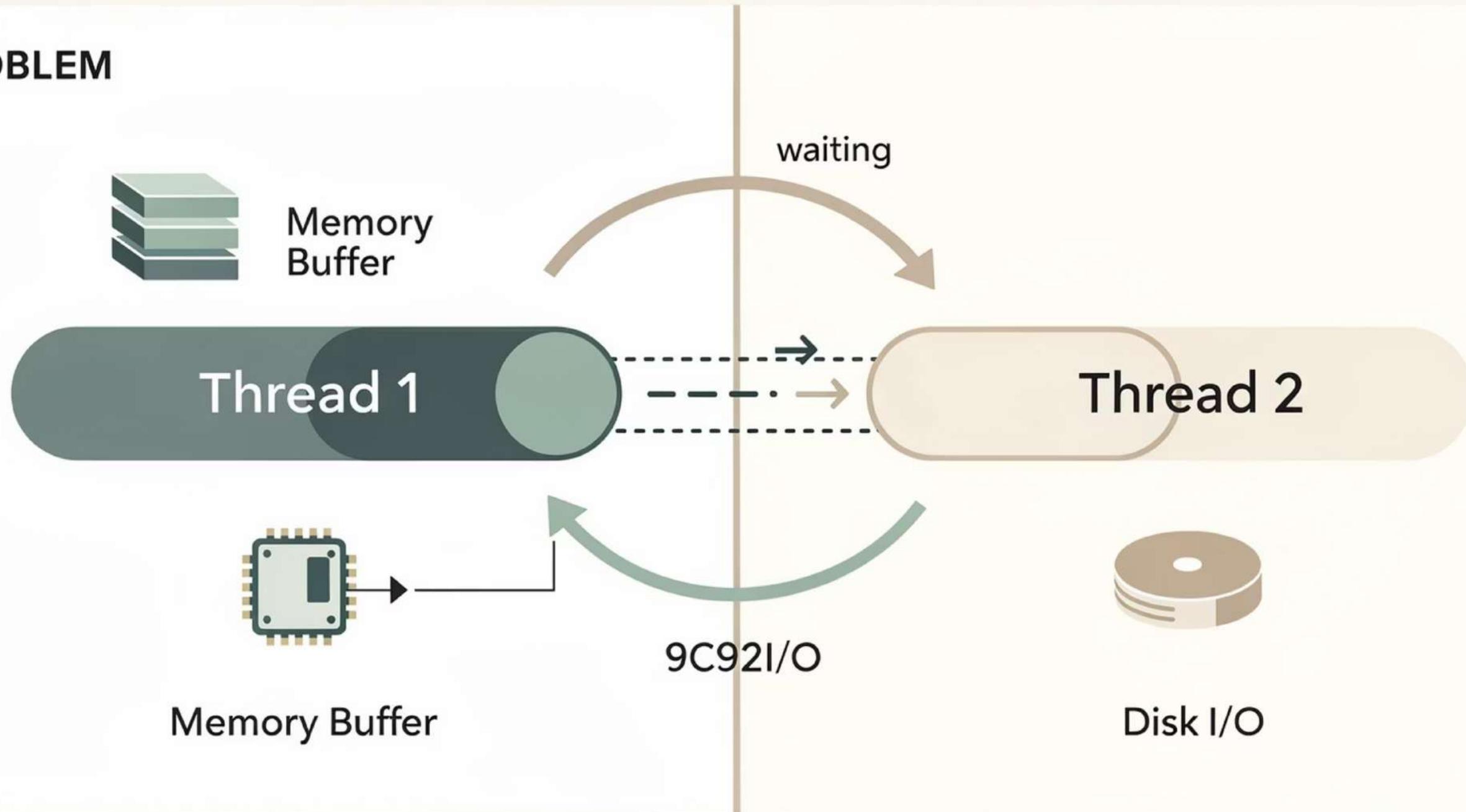
Three locks in the algorithm, per se, is a lipentha.

Customer Account

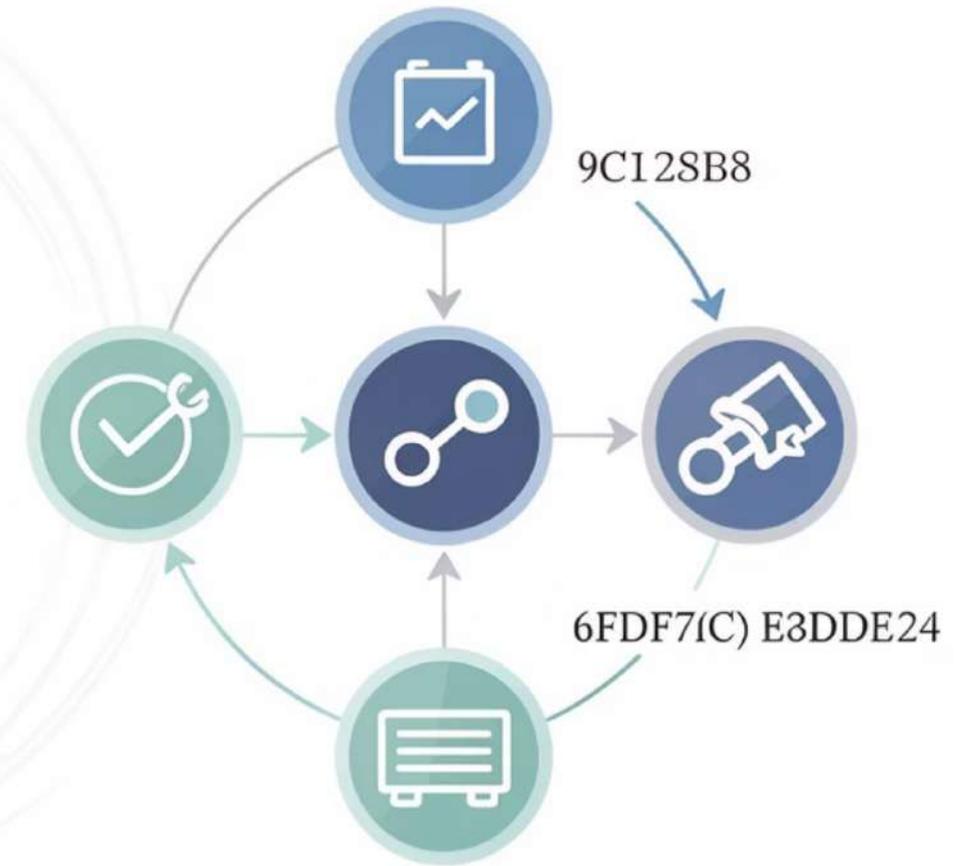
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Deadlock in Video Editing: A Case Study

PROBLEM



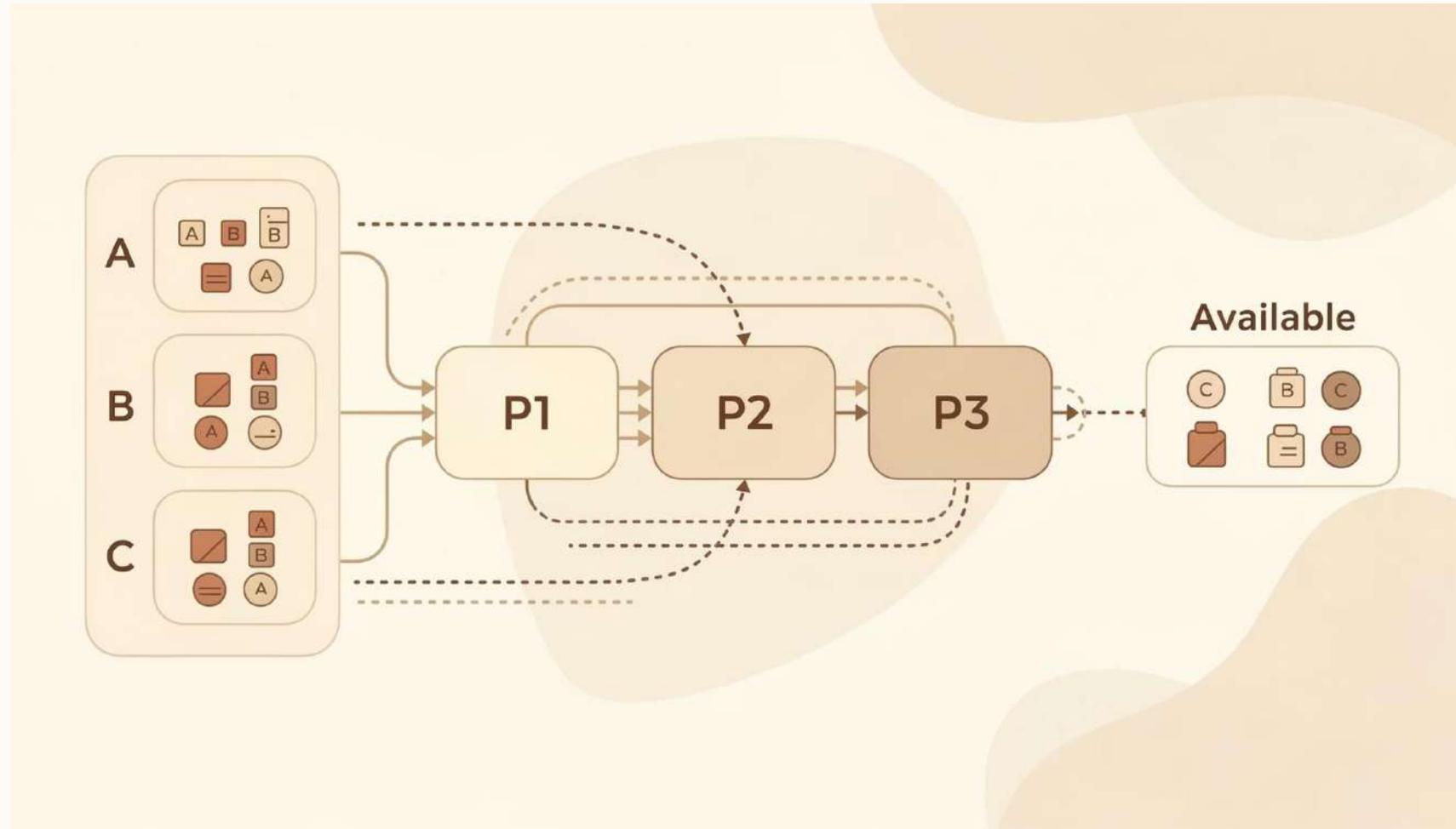
Deadlock Prevention & Avoidance



Prevention and Avoidance

The Resource Allocation Puzzle

A system has 3 processes (P1, P2, P3) and 3 resource types (A=10 instances, B=5 instances, C=7 instances).



Questions:

1. Is the system currently in a safe state?
2. If P2 requests (1, 0, 2), should the request be granted?
3. Explain your reasoning using the Banker's Algorithm

Puzzle Solution

01

Calculate Need Matrix

Need = Maximum - Allocated

P1: (7,4,3) | P2: (1,2,2) | P3: (6,0,0)

02

Find Safe Sequence

Available: (3,3,2) → P2 can finish (need ≤ available)

After P2: (5,3,2) → P1 can finish

After P1: (5,4,2) → P3 can finish

Safe sequence exists: P2 → P1 → P3

03

Evaluate P2's Request

Request (1,0,2). New available: (2,3,0)

P2's new need: (0,2,0). Check if any process can complete...

No safe sequence found → DENY request

❏ **Answer: Yes, currently safe. No, don't grant P2's request—it would create an unsafe state where no process can complete**



Thank You

