SNS COLLEGE OF TECHNOLOGY

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DEPARTMENT OF COMPUTER APPLICATIONS

COURSE

23CAE717
Cloud Computing

UNIT III

Cloud Infrastructure

TOPIC

Virtualization of CPU, Memory and I/O devices

Semester

II Semester / I MCA

CPU Virtualization

- Modern OS & processors permit multiple processes to run simultaneously
- All processors have at least two modes, user mode and supervisor mode
- Instructions running in supervisor mode are called privileged instructions
- Other instructions are unprivileged instructions
- VMware Workstation is a VM software suite for x86 and x86-64 computers
- KVM (Kernel-based Virtual Machine) is a Linux kernel virtualization infrastructure

CPU Virtualization

- unprivileged instructions of VMs run directly on the host machine
- critical instructions should be handled carefully
- Three categories of critical instructions: privileged instructions, control-sensitive instructions, and behavior-sensitive instructions
- Privileged instructions execute in a privileged mode and will be trapped if executed outside this mode.

CPU Virtualization

- Control-sensitive instructions attempt to change the configuration of resources used.
- Behavior-sensitive instructions have different behaviors depending on the configuration of resources, including the load/store operations over the virtual memory
- RISC CPU architectures can be naturally virtualized
- x86 CPU architectures are not primarily designed to support virtualization, because 10 sensitive instructions, are not privileged instructions

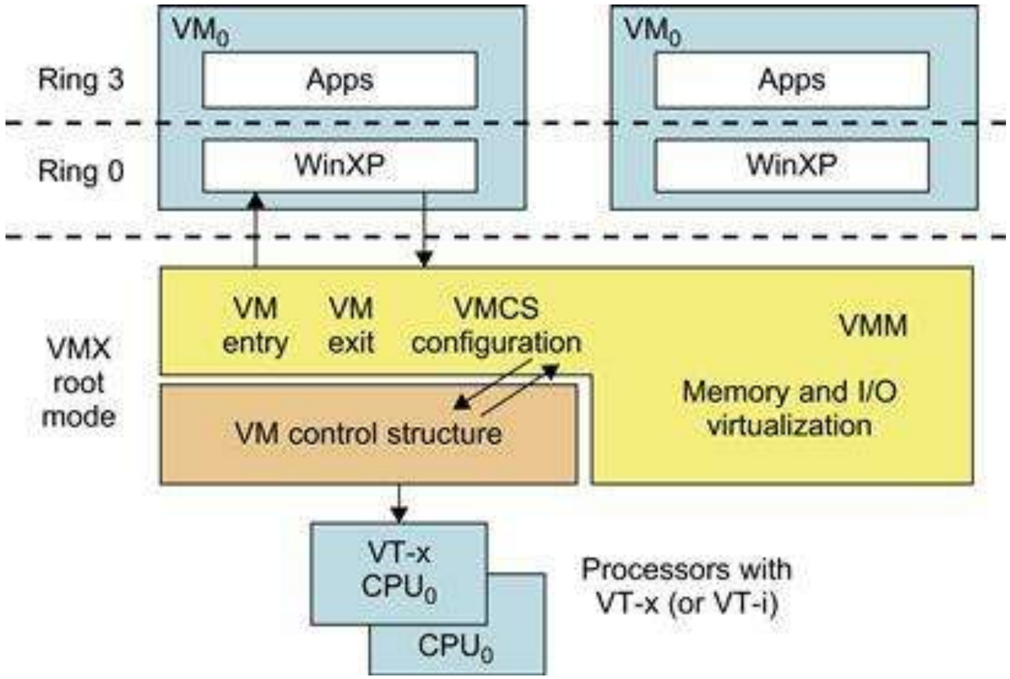
Hardware assisted CPU Virtualization

- Intel and AMD add an additional mode called privilege mode level (some people call it Ring-1) to x86 processors. Therefore, operating systems can still run at Ring 0 and the hypervisor can run at Ring -1.
- All the privileged and sensitive instructions are trapped in the hypervisor automatically.
- This technique removes the difficulty of implementing binary translation of full virtualization

CPU Virtualization

- A CPU architecture is virtualizable if it supports the ability to run the VM's privileged and unprivileged instructions in the CPU's user mode while the VMM runs in supervisor mode.
- When the privileged instructions including control- and behavior-sensitive instructions of a VM are executed, they are trapped in the VMM

Memory Virtualization

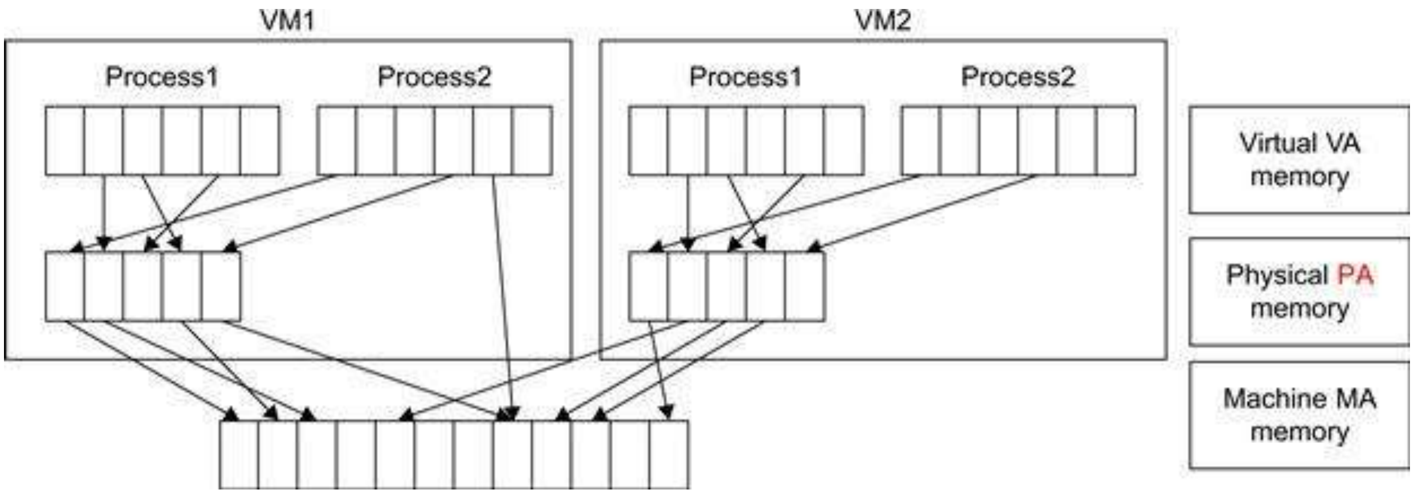


Intel hardware-assisted CPU virtualization

Memory Virtualization

- Virtual memory virtualization is similar to virtual memory supported by modern OS
- modern x86 CPUs include a memory management unit (MMU) and a translation lookaside buffer (TLB) to optimize virtual memory performance.
- Two-stage mapping process should be maintained by the guest OS and the VMM, respectively: virtual memory to physical memory and physical memory to machine memory

Memory Virtualization



Two-level memory mapping procedure

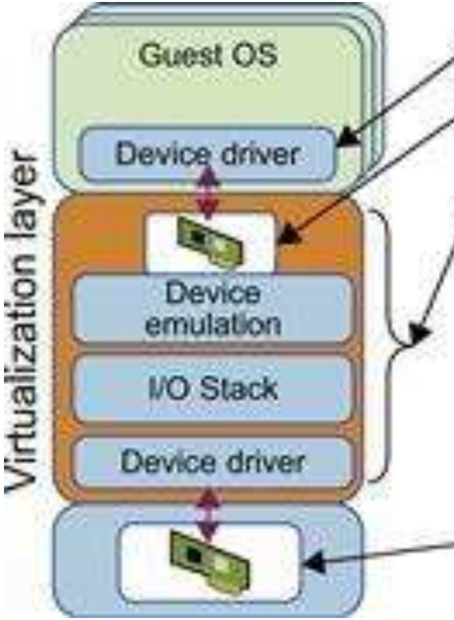
Memory Virtualization

- each page table of the guest OSes has a separate page table in the VMM corresponding to it, the VMM page table is called the shadow page table
- MMU already handles virtual-to-physical translations as defined by the OS
- VMware uses shadow page tables to perform virtual-memory-to-machine-memory address translation.
- Processors use TLB hardware to map the virtual memory directly tothe machine memory to avoid the two levels of translation on every access

I/O Virtualization

- involves managing the routing of I/O requests between virtual devices and the shared physical hardware
- three ways to implement I/O virtualization:
  - full device emulation
  - para-virtualization
  - direct I/O
- ***Full device emulation***: All the functions of a device or bus infrastructure, such as device enumeration, identification, interrupts, and DMA, are replicated in software

I/O Virtualization



- Guest device driver
- Virtual device
- Virtualization layer
  - emulates the virtual device
  - remaps guest and real I/O addresses
  - multiplexes and drives the physical device
  - I/O features. e.g., COW disks
- Real device
  - may be different from virtual device

I/O Virtualization

- ***Para-virtualization***: consisting of a frontend driver and a backend driver.
- The frontend driver is running in Domain U and the backend driver is running in Domain 0.
- The frontend driver manages the I/O requests of the guest OSes
- Backend driver is responsible for managing the real I/O devices and multiplexing the I/O data of different VM
- achieves better device performance than full device emulation

I/O Virtualization

- ***Direct I/O virtualization***: lets the VM access devices directly.
- It can achieve close-to- native performance without high CPU costs.
- self-virtualized I/O (SV-IO): All tasks associated with virtualizing an I/O device are encapsulated in SV-IO. It provides virtual devices and an associated access API to VMs and a management API to the VMM
- defines one virtual interface (VIF) for every kind of virtualized I/O device, such as virtual network interfaces, virtual block devices (disk), virtual camera devices,

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