



# Unit III – Dynamic Programming



- Dynamic Programming
  - Computing a Binomial Coefficient
  - Warshall's algorithm
  - **Floyd's algorithm**
  - **Optimal Binary Search Trees**
  - Knapsack Problem and Memory functions

# Floyd's algorithm

- Weighted connected graph – all pair shortest path
- Algorithm

```
ALGORITHM Floyd( $W[1..n, 1..n]$ )  
  //Implements Floyd's algorithm for the all-pairs shortest-paths problem  
  //Input: The weight matrix  $W$  of a graph with no negative-length cycle  
  //Output: The distance matrix of the shortest paths' lengths  
   $D \leftarrow W$  //is not necessary if  $W$  can be overwritten  
  for  $k \leftarrow 1$  to  $n$  do  
    for  $i \leftarrow 1$  to  $n$  do  
      for  $j \leftarrow 1$  to  $n$  do  
         $D[i, j] \leftarrow \min\{D[i, j], D[i, k] + D[k, j]\}$   
  return  $D$ 
```

- Time Complexity –  $O(n^3)$

# Optimal Binary Search Tree

Cost Matrix

$$C[i, i-1] = 0$$

$$C[i, i] = 0$$

$$C[i, j] = \text{formula}$$

Root Matrix

$$R[i, i] = i$$

$$R[i, j] = k \text{ (min)}$$

	0	1	2	3	4
1	<b>0</b>	<b>0.1</b>	0.4	1.1	1.7
2		<b>0</b>	<b>0.2</b>	0.8	1.4
3			<b>0</b>	<b>0.4</b>	1.0
4				<b>0</b>	<b>0.3</b>
5					<b>0</b>

	0	1	2	3	4
1		<b>1</b>	2	3	3
2			<b>2</b>	3	3
3				<b>3</b>	3
4					<b>4</b>
5					

CT

	e	1	2	3	4
1	0	0.1	0.4	1.1	1.7
2		0	0.2	0.8	1.4
3			0	0.4	1.0
4				0	0.3
5					0

RT

	0	1	2	3	4
1		1	2	3	3
2			2	3	3
3				3	3
4					4
5					

$C[i, i-1] = 0$

$C[i, i] = p_i$

$C[i, j] = \text{Formula}$

$R[i, i] = i$

$R[i, j] = k \text{ (minimum)}$

Table  $\rightarrow$  Upper right corner  $\Rightarrow$  Tree Construction

$(1, 4) \quad k=3$

$i, j$

```

graph TD
    C((C)) --> B((B))
    C --> D((D))
    B --> A((A))
            
```

$A, B, C, D \Rightarrow 0.1, 0.2, 0.4, 0.3$

$1 \quad 2 \quad 3 \quad 4$

```

graph TD
    R((R)) --> 2((2))
    R --> 4((4))
    2 --> 1((1))
    4 --> 3((3))
            
```

$(i, k-1) \quad L$   
 $(1, 2)$

$k+1, j$   
 $(4, 4)$