Ford-Fulkerson Algorithm

Initially, the flow of value is 0. Find some augmenting Path p and increase flow f on each edge of p by residual Capacity $c_f(p)$. When no augmenting path exists, flow f is a maximum flow.

FORD-FULKERSON METHOD (G, s, t)

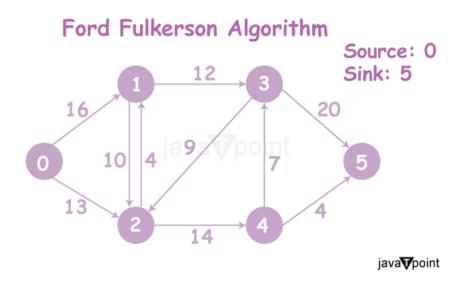
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1. Initialize flow f to 0
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- 2. while there exists an augmenting path $\ensuremath{\mathsf{p}}$
- 3. do argument flow f along $\ensuremath{\mathsf{p}}$
- 4. Return f

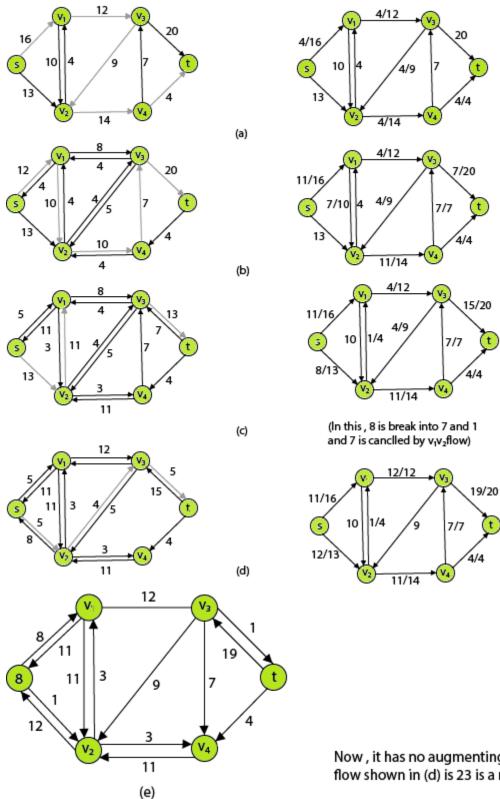
FORD-FULKERSON (G, s, t)

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for each edge (u, v) ∈ E [G]
do f [u, v] ← 0
f [u, v] ← 0
while there exists a path p from s to t in the residual network G<sub>f</sub>.
do c<sub>f</sub> (p)←min?{ C<sub>f</sub> (u,v): (u,v) is on p}
for each edge (u, v) in p
do f [u, v] ← f [u, v] + c<sub>f</sub> (p)
f [u, v] ←-f[u,v]
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Example: Each Directed Edge is labeled with capacity. Use the Ford-Fulkerson algorithm to find the maximum flow.



Solution: The left side of each part shows the residual network Gf with a shaded augmenting path p,and the right side of each part shows the net flow f.



Now, it has no augmenting paths. So, the maximum flow shown in (d) is 23 is a maximum flow .