



SNS COLLEGE OF TECHNOLOGY

An Autonomous Institution

Coimbatore-35



Accredited by NBA – AICTE and Accredited by NAAC – UGC with 'A+' Grade
Approved by AICTE, New Delhi & Affiliated to Anna University, Chennai

DEPARTMENT OF ELECTRONICS & COMMUNICATION ENGINEERING

19ECB212 – DIGITAL SIGNAL PROCESSING

II YEAR/ IV SEMESTER

UNIT 4 – FINITE WORD LENGTH EFFECTS

TOPIC – INTRODUCTION TO DSP PROCESSOR



INTRODUCTION



In Digital Signal Processors (DSPs) are microprocessors with the characteristics:

- a) Real-time digital signal processing capabilities. DSPs typically have to process data in real time, i.e., the correctness of the operation depends heavily on the time when the data processing is completed
- b) High throughput. DSPs can sustain processing of high-speed streaming data, such as audio and multimedia data processing
- c) Deterministic operation. The execution time of DSP programs can be foreseen accurately, thus guaranteeing a repeatable, desired performance



INTRODUCTION



d) Re-programmability by software. Different system behavior might be obtained by re-coding the algorithm executed by the DSP instead of by hardware modifications

- DSPs appeared on the market in the early 1980s. Over the last 15 years they have been the key enabling technology for many electronics products in fields such as commn systems, multimedia, automotive, instrumentation & military
- The DSP implements the audio and encode functions. Additional tasks carried out are file management, user interface control, and post-processing algorithms such as equalization and bass management



SPECIFIC APPLICATIONS OF DSP



Field		Application
Communication	Broadband	Video conferencing / phone
		Voice / multimedia over IP
		Digital media gateways (VOD)
	Wireless	Satellite phone
		Base station
	Consumer	Security
Video surveillance		
Entertainment		Digital still /video camera
		Digital radio
		Portable media player / entertainment console
Toys		Interactive toys
		Video game console



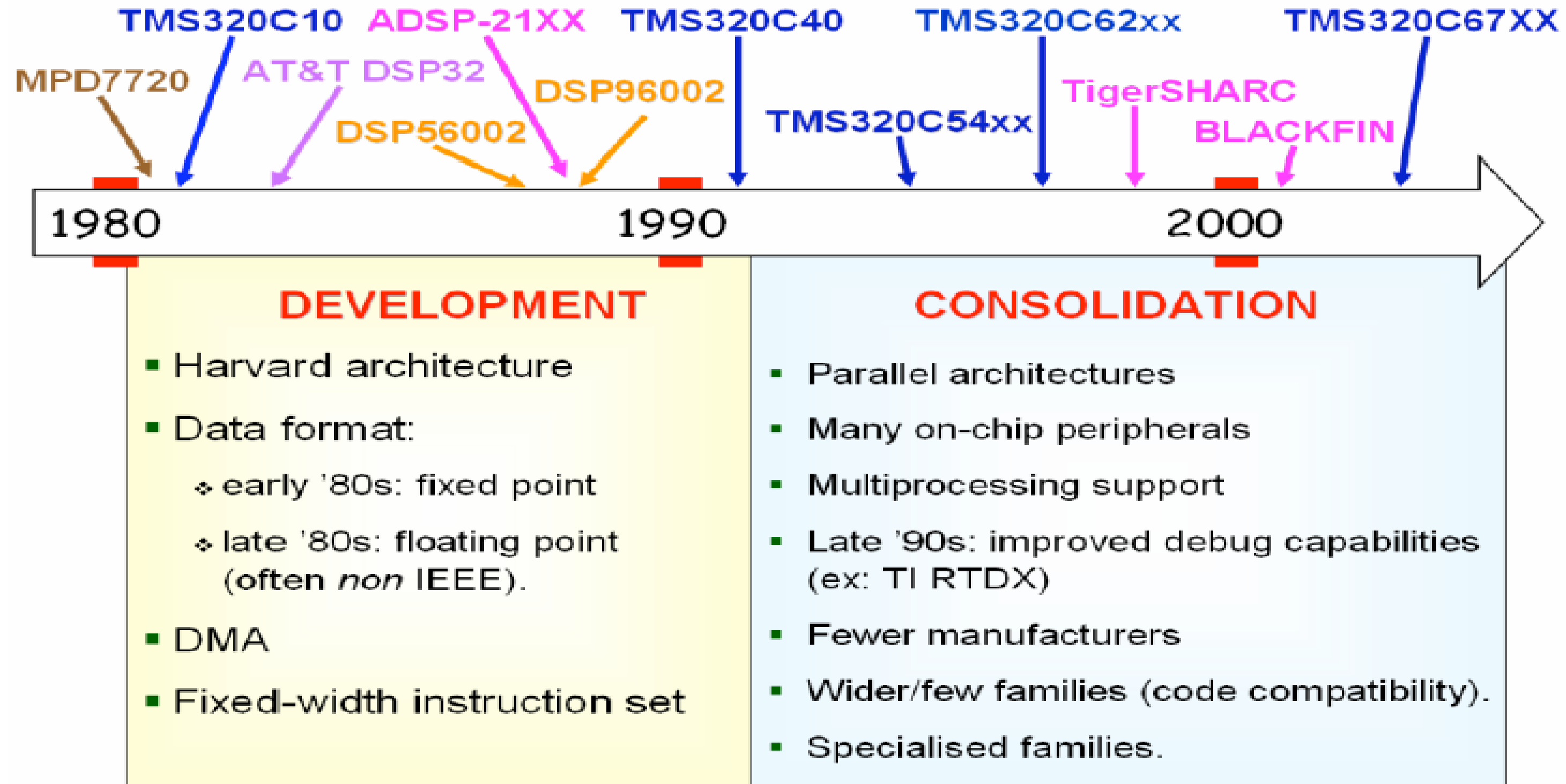
SPECIFIC APPLICATIONS OF DSP



Consumer	Security	Biometrics
		Video surveillance
	Entertainment	Digital still /video camera
		Digital radio
		Portable media player / entertainment console
	Toys	Interactive toys
		Video game console
	Industrial and entertainment	Medical
Ultrasound		
X-ray		
Point of sale		Scanner
		Vending machine
Industrial		Factory automation
		Industrial / machine / motor control
		Vision system



EVOLUTION OF DSP FEATURES FROM THEIR EARLY DAYS UNTIL NOW





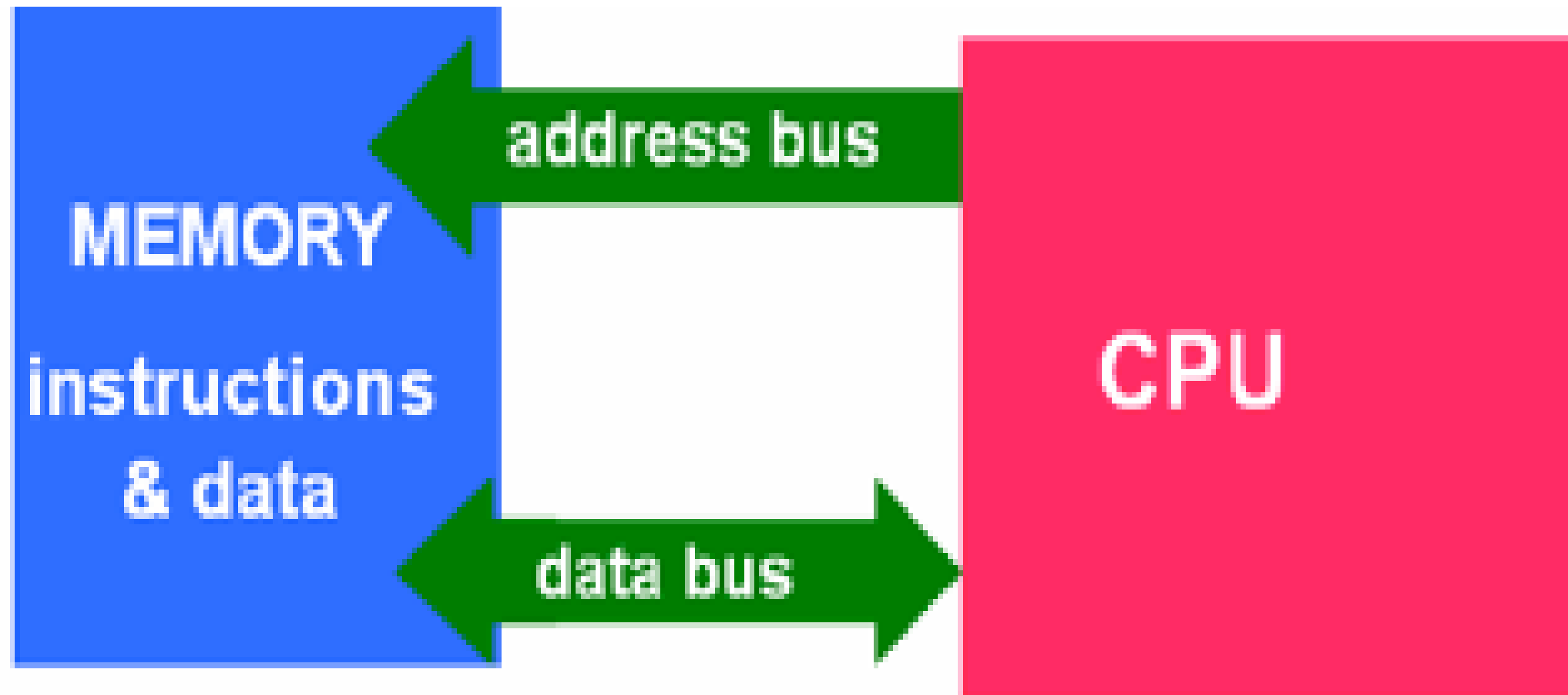
MAIN REQUIREMENTS AND HARDWARE IMPLEMENTATIONS FOR REAL-TIME DSP



Processing requirements	Hardware implementations satisfying the requirement
Fast data access	<ul style="list-style-type: none">• High-bandwidth memory architectures• Specialized addressing modes• Direct Memory Access (DMA)
Fast computation	<ul style="list-style-type: none">• MAC-centred• Pipelining• Parallel architectures (VLIW, SIMD)
Numerical fidelity	<ul style="list-style-type: none">• Wide accumulator registers, guard bits, etc.
Fast execution control	<ul style="list-style-type: none">• Hardware-assisted, zero-overhead loops, shadow registers, etc.

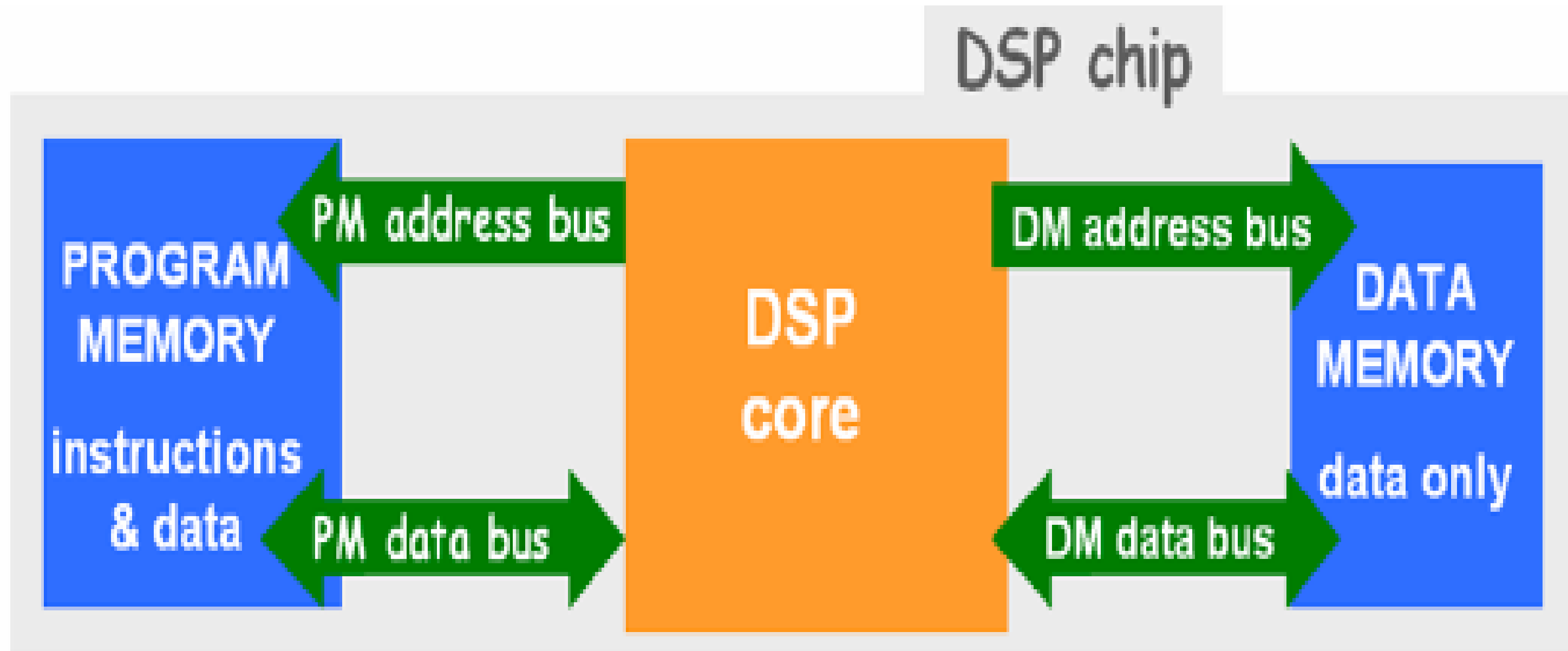


VON NEUMANN ARCHITECTURE



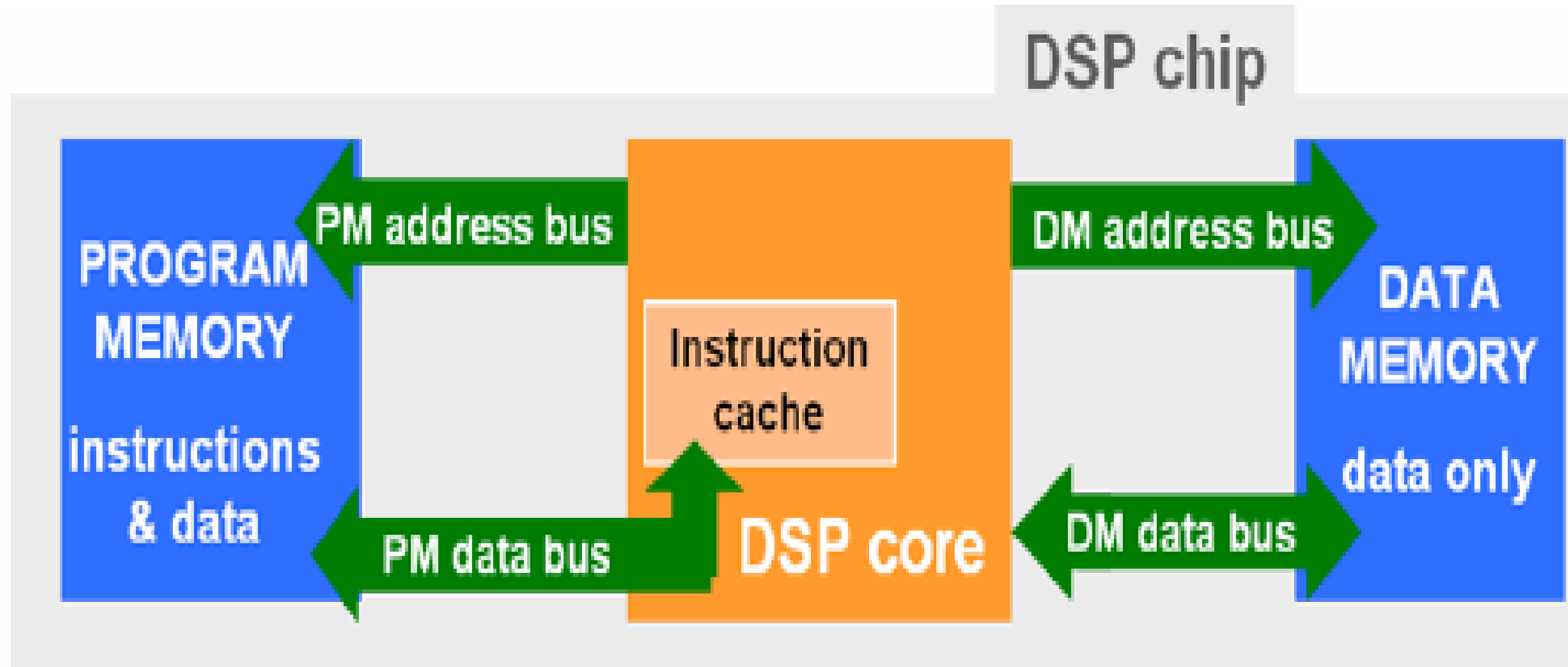


HARVARD ARCHITECTURE





SUPER HARVARD ARCHITECTURE





VON NEUMANN & HARVARD ARCHITECTURE



- Von Neumann Architecture consists of a single block of memory, containing both data and program instructions, and of a single bus (called data bus) to transfer data and instructions from/to the CPU
- In the Harvard architecture, there are separate memories for data and program instructions, and two separate buses connect them to the DSP core. This allows fetching program instructions and data at the same time, thus providing better performance at the price of an increased hardware complexity and cost
- The Harvard architecture can be improved by adding to the DSP core a small bank of fast memory, called 'instruction cache', and allowing data to be stored in the program memory



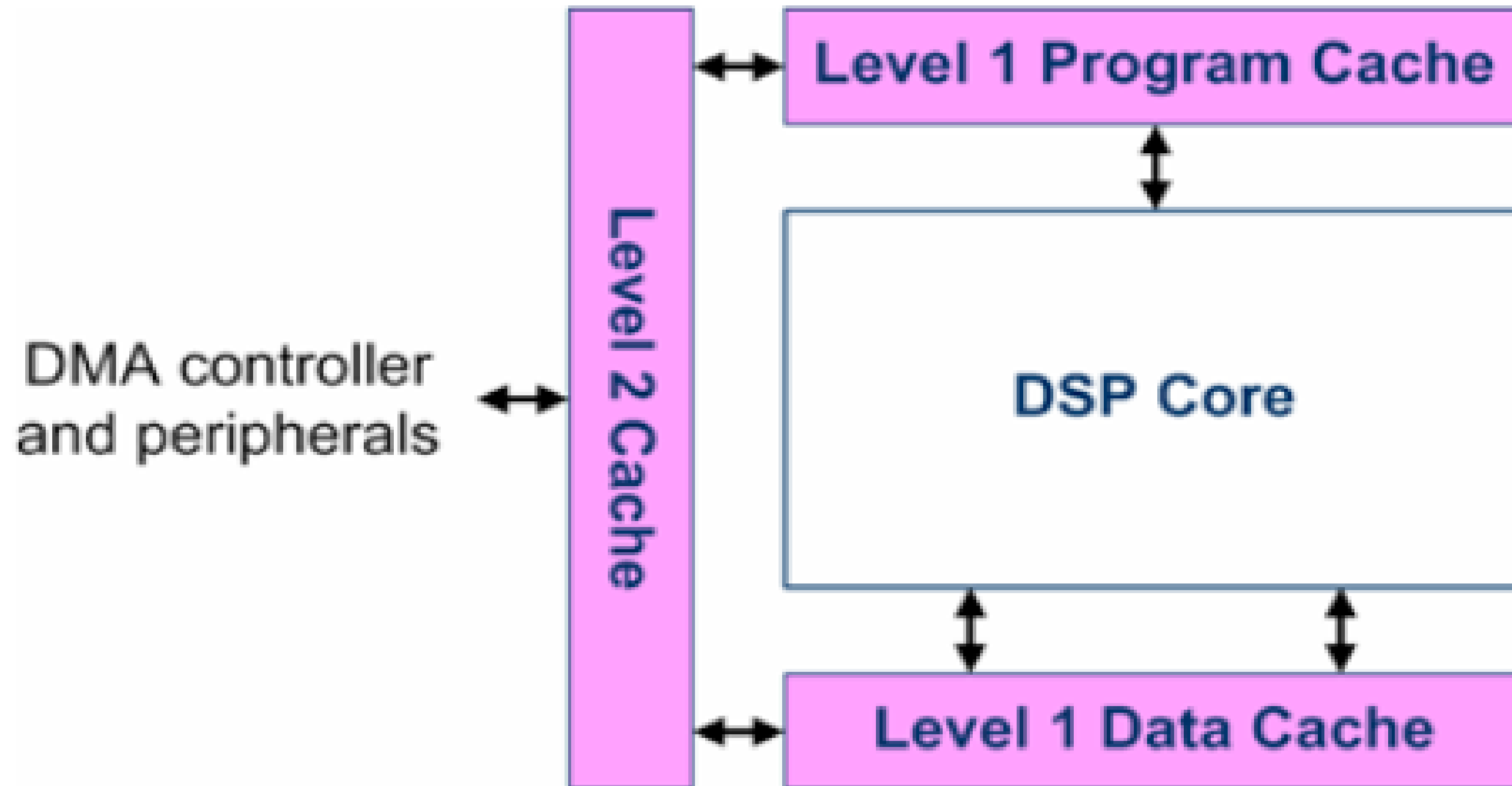
CACHE ARCHITECTURE



- The last-executed program instructions are relocated at run time in the instruction cache. The instructions are copied to the instruction cache the first time the DSP executes the loop.
- Cache architecture for TI TMS320C67xx DSP, including both program and data cache. There are two levels of cache, called Level 1 (L1) and Level 2 (L2). The L1 cache comprises 8 kbyte of memory divided into 4 kbyte of program cache and 4 kbyte of data cache. The L2 cache comprises 256 kbyte of memory divided into 192 kbyte mapped-SRAM memory and 64 kbyte dual cache memory



TI DSP TMS320C67XX FAMILY TWO-LEVEL CACHE ARCHITECTURE

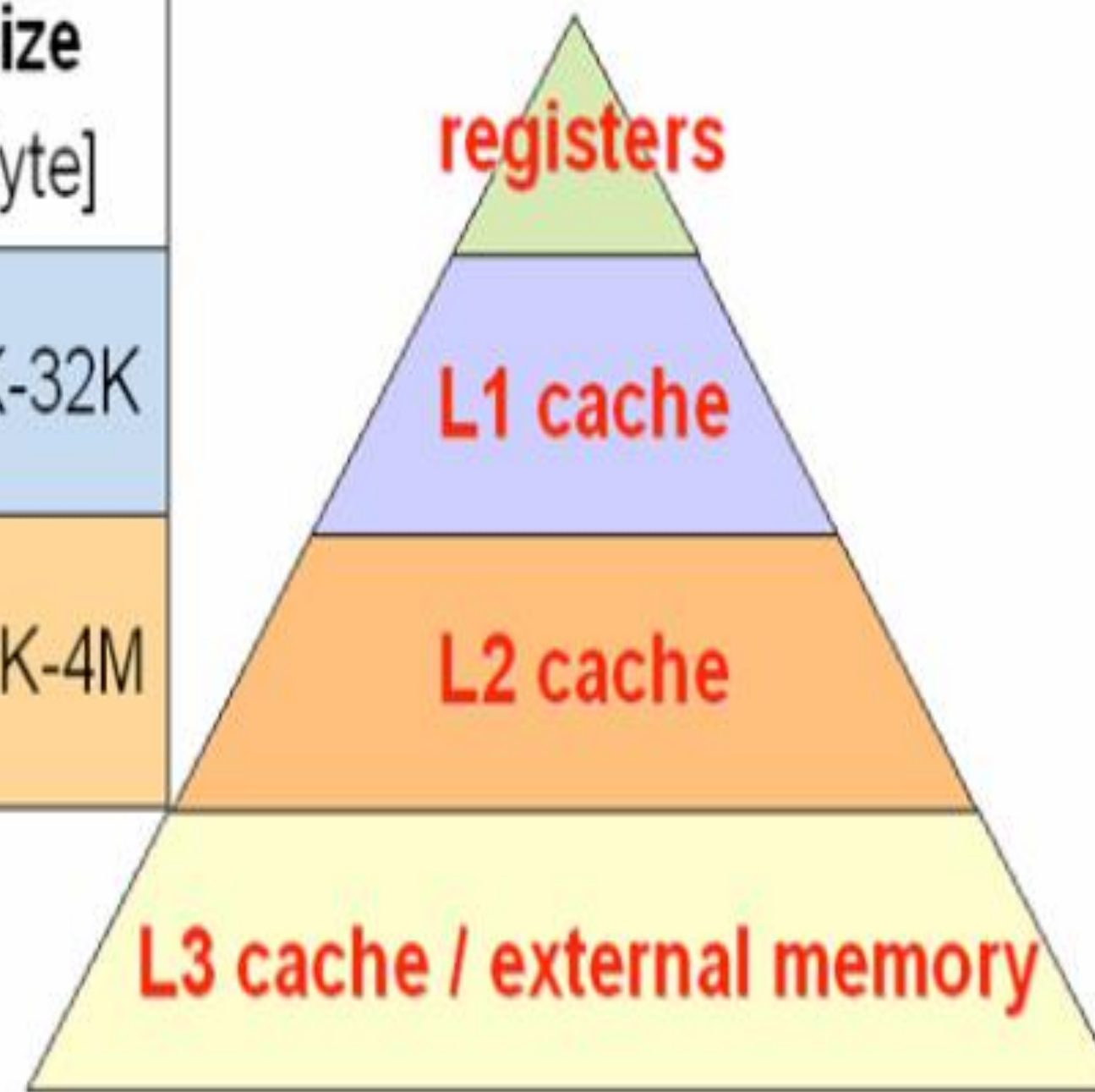




DSP HIERARCHICAL MEMORY ARCHITECTURE



Access [ns]	Hardware implementation	Size [Byte]
1	~5 transistor /cell	16K-32K
5-10	~2 transistor /cell	512K-4M
10-50		



smaller & faster



larger & slower



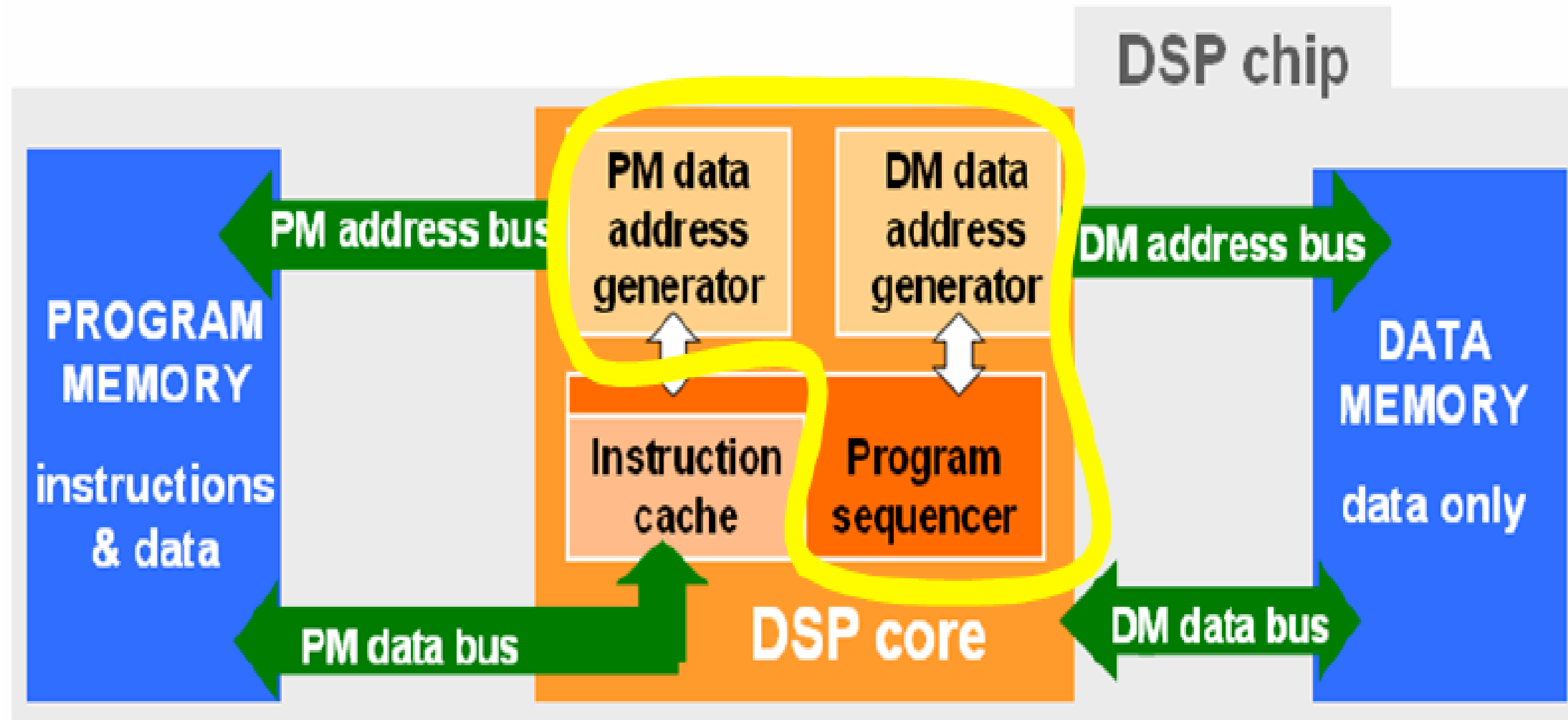
DSP HIERARCHICAL MEMORY ARCHITECTURE



- Hierarchical memory allows one to take advantage of both the speed and the capacity of different memory types. Registers are banks of very fast internal memory, typically with single-cycle access time
- The L1 cache is typically high-speed static RAM made of five or six transistors. The amount of L1 cache available thus depends directly on the available chip space. A L2 cache needs typically a smaller number of transistors hence can be present in higher quantities inside the DSPs
- A missing cache hit happens when the data or the instructions needed by the DSP are not stored in cache memory, hence they have to be fetched from a slower memory with an execution speed penalty

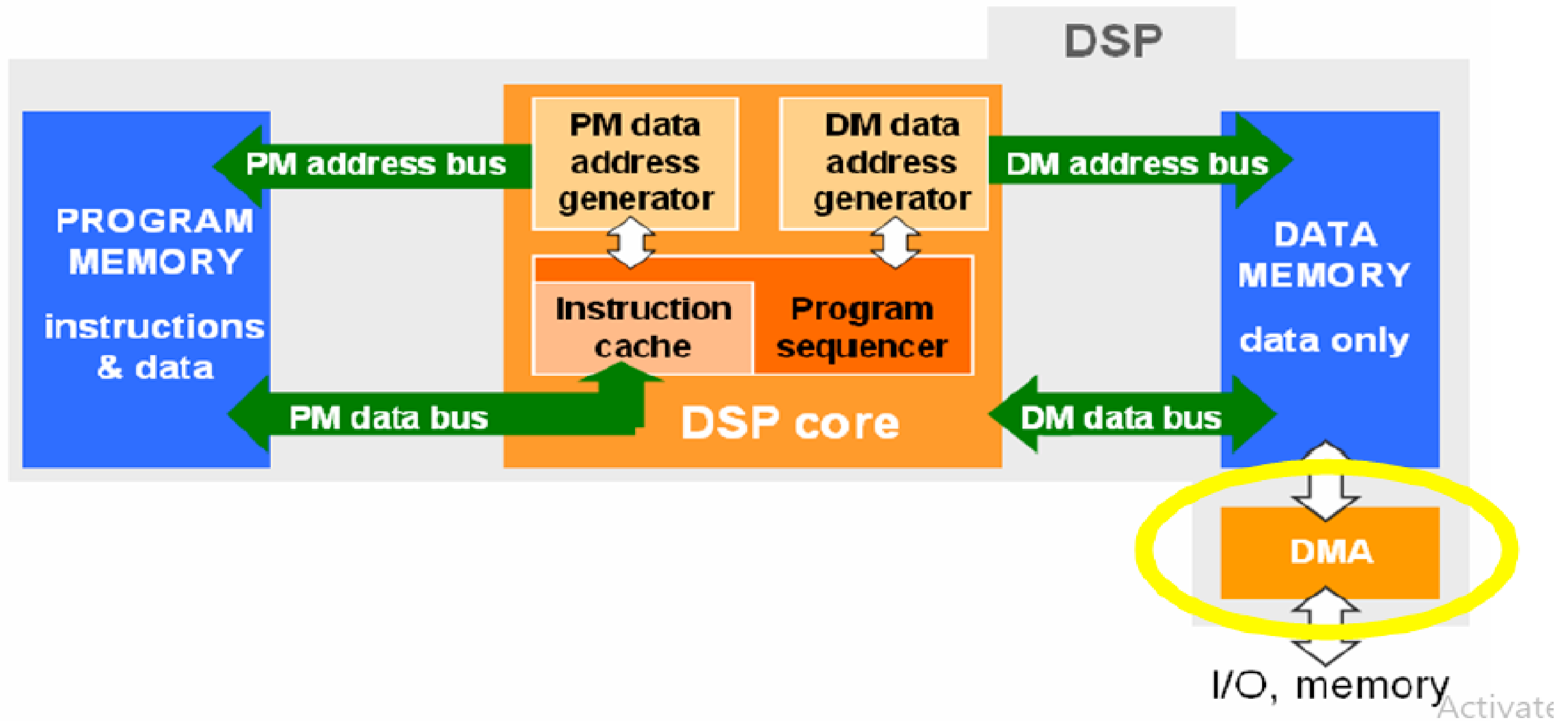


PROGRAM SEQUENCER AND ADDRESS GENERATOR UNITS LOCATION WITHIN A GENERIC DSP CORE ARCHITECTURE





DMA CONTROLLER





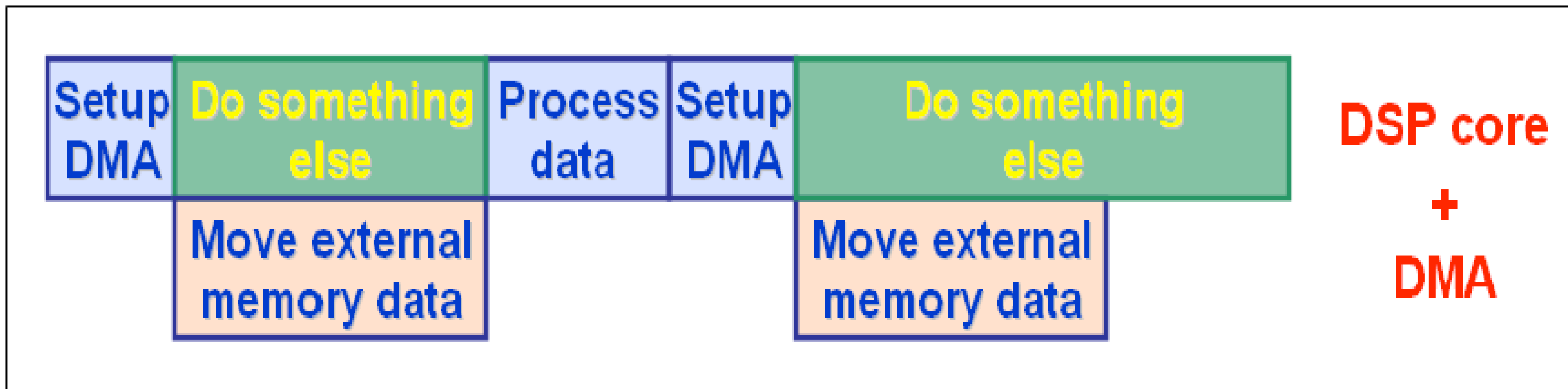
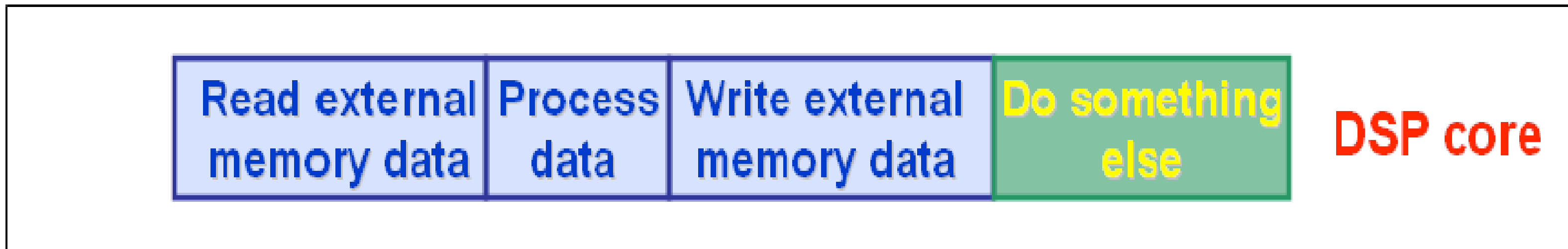
DMA CONTROLLER



- The DMA controller is a second processor working in parallel with the DSP core and dedicated to transferring information between two memory areas or between peripherals and memory
- A DMA coprocessor can transfer data as well as program instructions, the latter transfer corresponding typically to the case of code overlay, i.e., of code stored in an external memory and moved to an internal memory (for instance L1) when needed
- Multiple and independent DMA channels are also available for greater flexibility
- Bus arbitration between the DMA and the DSP core is needed to avoid colliding memory accesses when the DMA and the DSP core share the same bus to access peripherals and/or memories

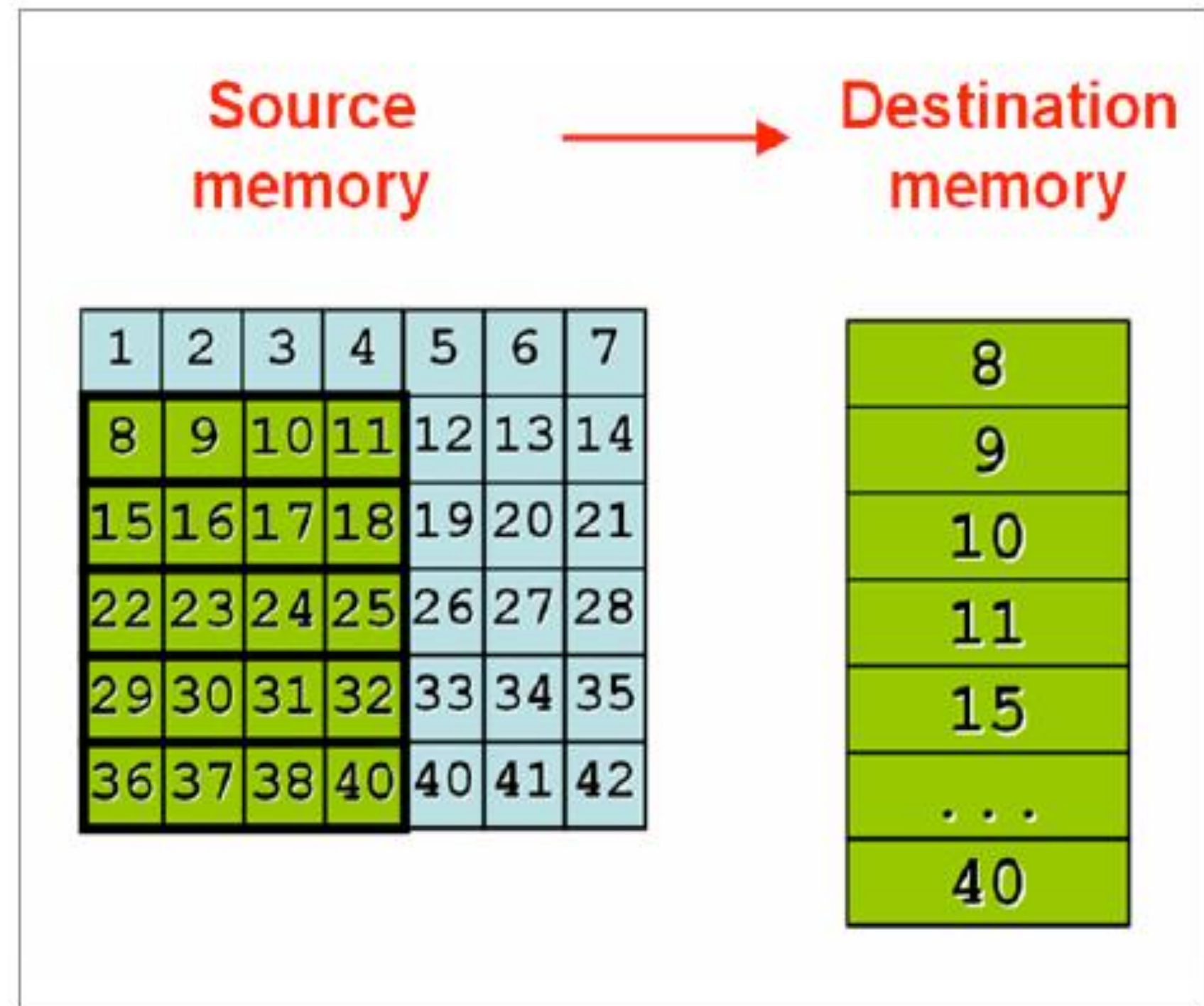
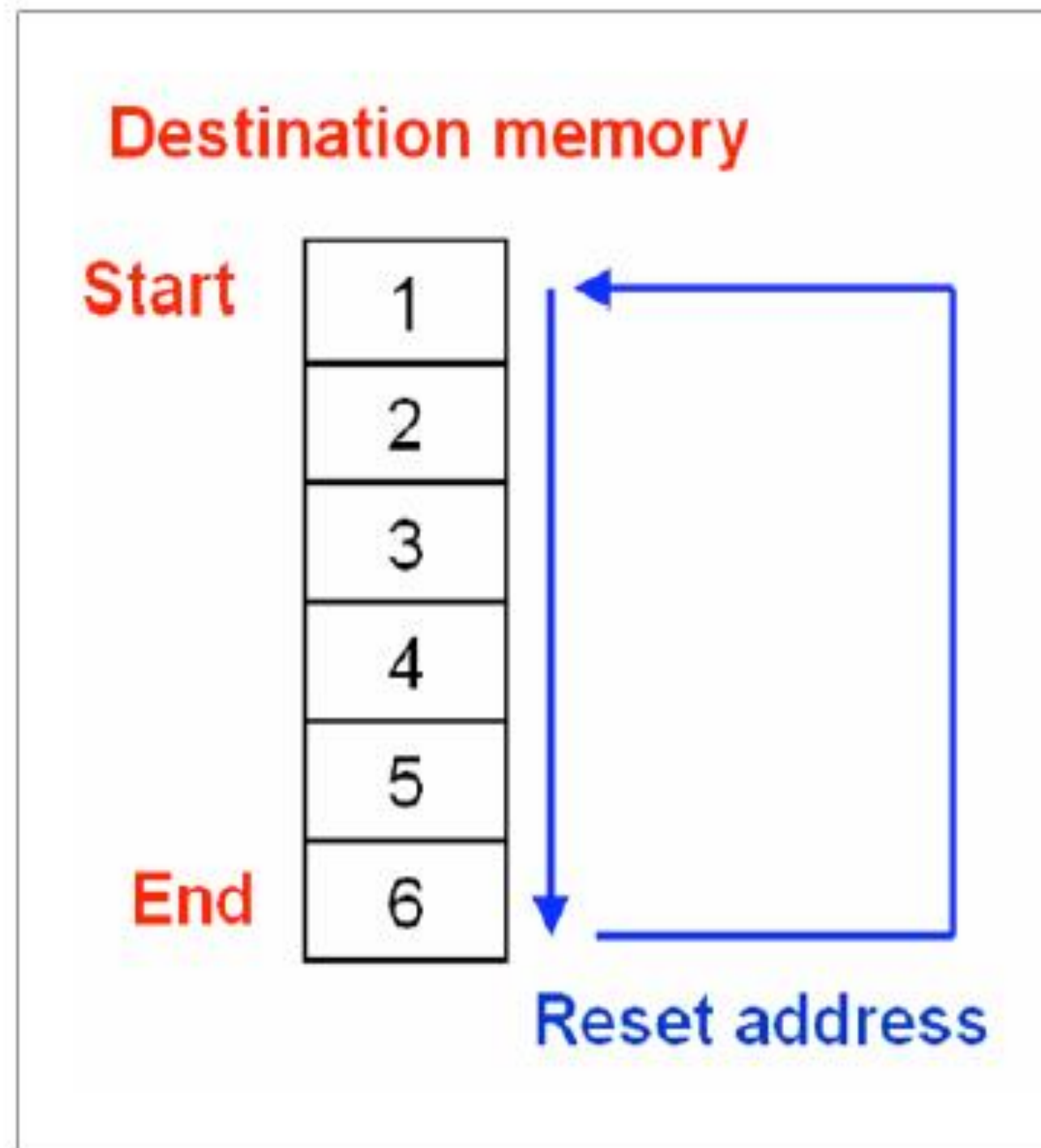


READ-PROCESS-WRITE DATA WHEN THE DSP CORE ONLY IS PRESENT; (B) SAME ACTIVITY WHEN THE DMA TAKES CARE OF DATA TRANSFERS





DMA TRANSFER CONFIGURATIONS. (A): CHAINED DMA TRANSFER; (B): MULTIDIMENSIONAL DATA TRANSFER





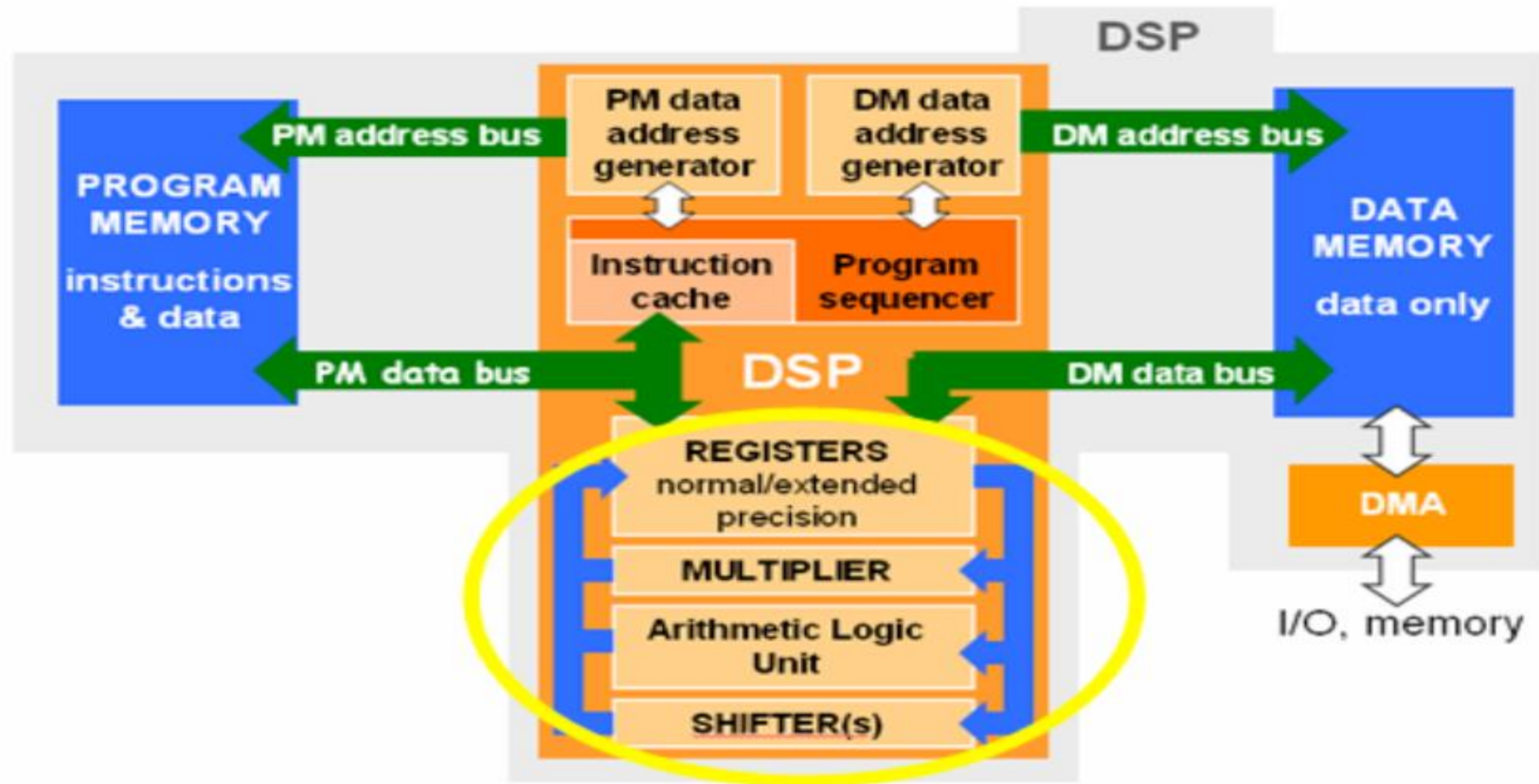
DSP PROCESSING ARITHMETIC BLOCKS



- The basic DSP arithmetic processing blocks are a) many registers; b) one or more multipliers; c) one or more Arithmetic Logic Units (ALUs); d) one or more shifters.
- a) Registers:** these are banks of very fast memory used to store intermediate data processing. Very often they are wider than the DSP normal word width, so as to provide a higher resolution during the processing
- b) Multiplier:** it can carry out single-cycle multiplications and very often it includes very wide accumulator registers to reduce round-off or truncation errors
- c) ALU:** it carries out arithmetic and logical operations.
- d) Shifters:** it shifts the input value by one or more bits, left or right. The shifter is called a barrel shifter and is useful in the implementation of floating point add and subtract operations



DSP PROCESSING ARITHMETIC BLOCKS





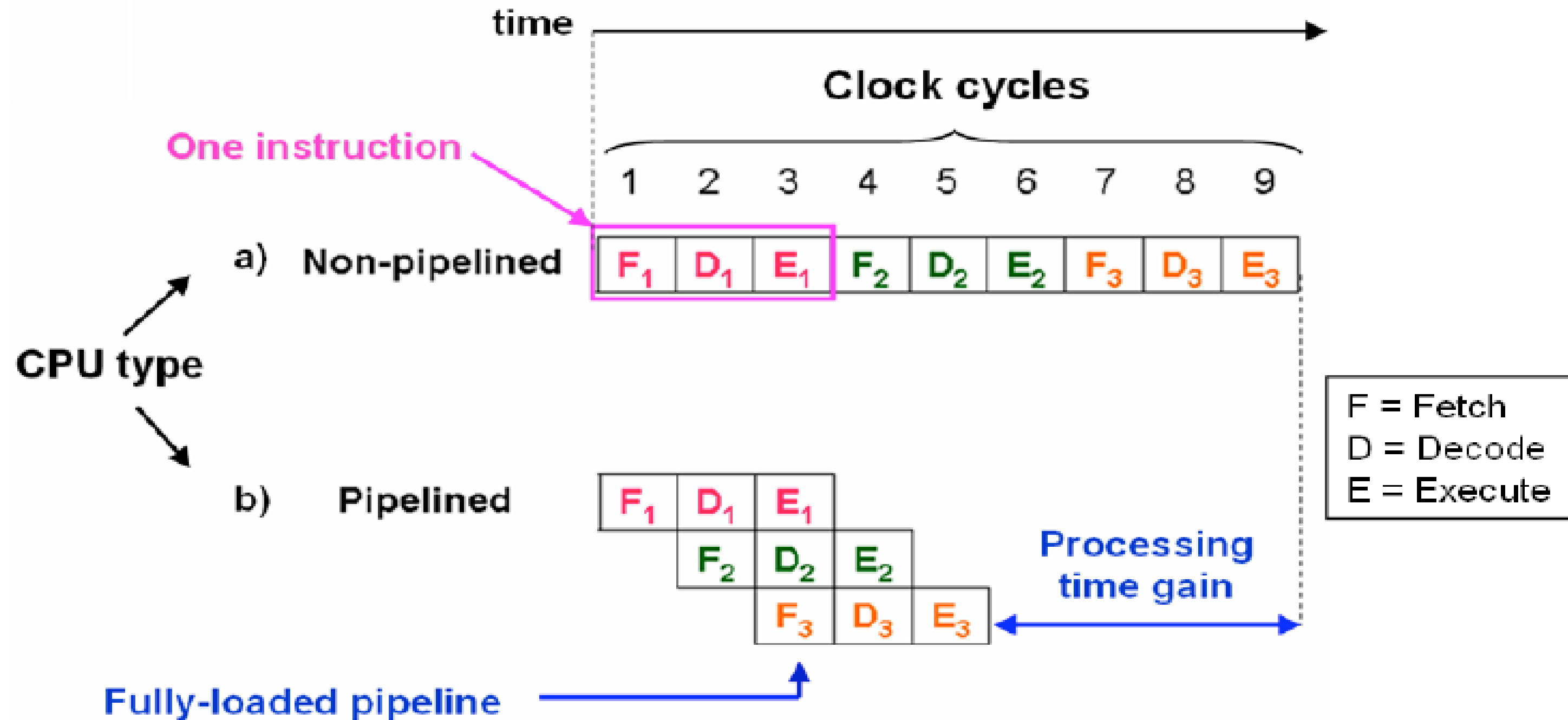
PIPELINING STAGES



Basic pipelining stages	Action
Fetch	<ul style="list-style-type: none">• Generate program fetch address• Read op-code
Decode	<ul style="list-style-type: none">• Route op-code to functional unit• Decode instruction• Read operands
Execute	<ul style="list-style-type: none">• Execute instruction• Write results back to registers



INSTRUCTION EXECUTION AND PROCESSING TIME GAIN OF A PIPELINED CPU WITH RESPECT TO A NON-PIPELINED ONE





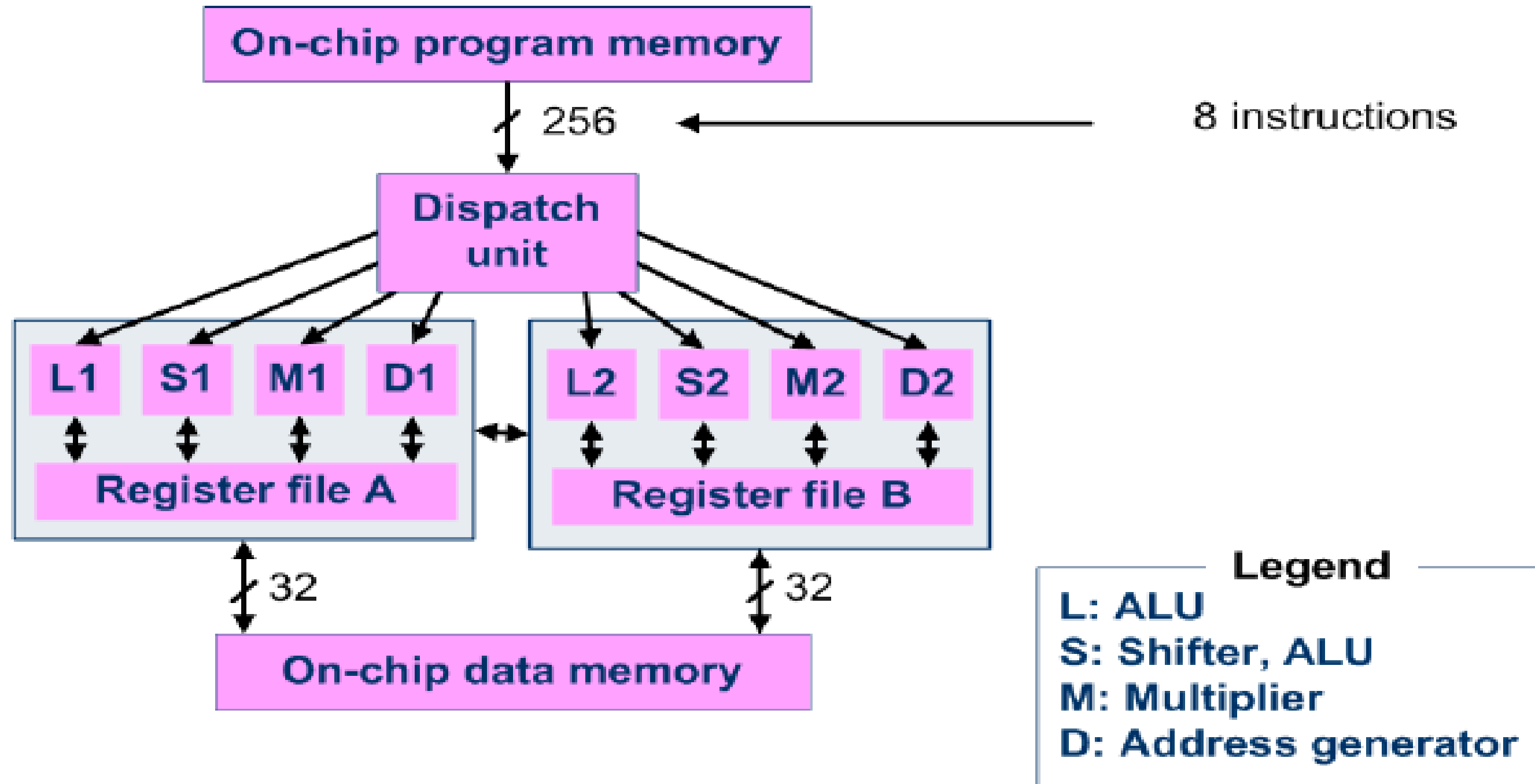
VLIW ARCHITECTURE



- VLIW architectures are based upon instruction level parallelism, i.e., many instructions are issued at the same time and are executed in parallel by multiple execution units. As a consequence, DSPs based on this architecture are also called 'multi-issue' DSP
- VLIW architecture: eight, 32-bit instructions are packed together in a 256-bit wide instruction which is fed to eight separate execution units. Characteristics of VLIW architectures include simple and regular instruction sets. Instruction scheduling is done at compile-time and not at run-time so as to guarantee a deterministic behavior



TI TMS320C6XXX FAMILY VLIW ARCHITECTURE





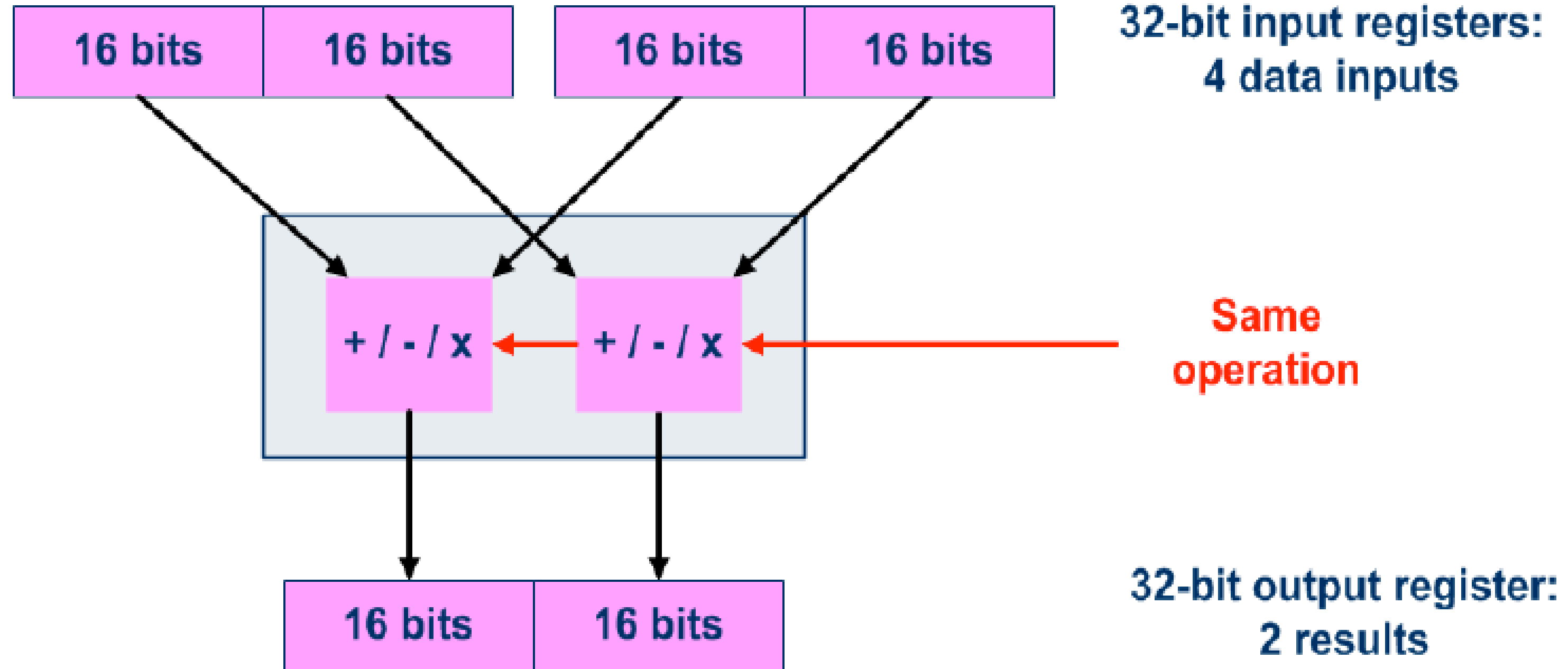
SIMD ARCHITECTURE



- SIMD architectures are based on data-level parallelism, i.e., only one instruction is issued at a time but the same operation specified by the instruction is performed on multiple data sets
- DSP based upon the SIMD architecture: two 32-bit input registers provide four, 16-bit each, data inputs. They are processed in parallel by two separate execution units that carry out the same operation. The two, 16-bit data outputs are packed into a 32-bit register
- Typical SIMD architecture can support multiple data width and is most effective on algorithms that require the processing of large data chunks. The SIMD operation mode can be switched ON or OFF, for instance in the ADI SHARC DSP



SIMPLIFIED SCHEMATICS FOR ADI SHARC DSP - SIMD ARCHITECTURE





ASSESSMENT



1. What are the characteristics of DSP?
2. List the main requirements and hardware implementations of real time DSP.
3. ----- Architecture consists of a single block of memory, containing both data and program instructions
4. What is meant by Harvard Architecture?
5. The basic DSP arithmetic processing blocks are -----, -----, ----- and -

6. Define VLIW Architecture.
7. List the basic pipelining stages.



THANK YOU